

# (A.I.)

Intelligent:  $\rightarrow$

- Who has sufficient knowledge. &
- Has Ability of reasoning using his knowledge

Diff. b/w Data & knowledge:-

- | Data  | Knowledge   |
|---|---|
| ① Data is stored in the form of tables                        | ① Knowledge is stored in many forms                                   |
| ② In <del>the</del> database we just search the data exactly. | ② In knowledge representation some methods are provided for reasoning |

Knowledge Representation  $\rightarrow$

1. Proposition logic  $\rightarrow$   
(zero order logic)

- if it rains <sup>then</sup> road will be wet.
- It is raining.

Here It rains : R } These are propositions  
 road is wet : W }

Propositions represent sentences.

Knowledge  $\left\{ \begin{array}{l} 1. R \rightarrow W \rightarrow \text{means If } R \text{ then } W \\ 2. R \end{array} \right.$

Make propositional logic  $\rightarrow$

1. If it is hot and plant is open then plant will be dead.

2. It is hot.

3. Plant is open.

(i) It is hot : H

(ii) Plant is open : P

(iii) Plant will be dead : D

1.  $H \wedge P \rightarrow D$

Means If  $H \wedge P$  then D

2. H

3. P

Definition  $\rightarrow$  Propositional logic formula  $\rightarrow$

1. F, T & F are Propositional Logic Form (PLF)

2. A, B, C etc are PLF where A, B, C are propositions and each proposition represents sentence/statements.

3. Item discussed in point 1 & 2 combine with following operators will

also be a propositional logic Formulae  
Operators  $\rightarrow$

1.  $\neg$  - Negation (Unary operator)  
 $\neg A$   
only 1 op required

2.  $\wedge$  - Conjunction / AND

3.  $\vee$  - Disjunction / OR

} Binary operators

4.  $\rightarrow$  - Implies

5.  $\leftrightarrow$  - Equivalence / Equal to

Precedence decreases downward.

### Mathematical Laws $\rightarrow$

Laws of propositional & predicate logic  $\rightarrow$

1. Commutative Laws

$$A \vee B = B \vee A$$

$$A \wedge B = B \wedge A$$

$$A \leftrightarrow B = B \leftrightarrow A$$

$$A \rightarrow B \neq B \rightarrow A \text{ (Implies is not commutative)}$$

2

2. Associative Laws  $\rightarrow$

$$A \vee (B \vee C) = (A \vee B) \vee C$$

$$A \wedge (B \wedge C) = (A \wedge B) \wedge C$$

3. Distributive Laws  $\rightarrow$

$$A \vee (B \wedge C) = (A \vee B) \wedge (A \vee C)$$

$$A \wedge (B \vee C) = (A \wedge B) \vee (A \wedge C)$$

4. De Morgan's Laws  $\rightarrow$

$$\neg (A \wedge B) = \neg A \vee \neg B$$

$$\neg (A \vee B) = \neg A \wedge \neg B$$

5. Law of Negation  $\rightarrow$

$$\neg(\neg A) = A$$

6. Law of Excluded Middle  $\rightarrow$

$$A \vee \neg A = T \text{ (True)}$$

{ Because  $T \vee F = F \vee T = T$  }  
This is tautology

7. Law of Contradiction  $\rightarrow$

$$A \wedge \neg A = F$$

This is Contradiction.

8. Law of Implication  $\rightarrow$

$$A \rightarrow B = \neg A \vee B$$

9. Law of equality  $\Rightarrow$

$$A \leftrightarrow B = (A \rightarrow B) \wedge (B \rightarrow A)$$

10. Laws of OR-simplification

$$A \vee A = A$$

$$A \vee T = T$$

$$A \vee F = A$$

$$A \vee (A \wedge B) = A$$

11. Laws of AND Simplification

$$A \wedge A = A$$

$$A \wedge T = A$$

$$A \wedge F = F$$

$$A \wedge (A \vee B) = A$$

Q. Simplify the following logical expr<sup>n</sup>:

1.  $\neg(b \wedge \neg b)$

2.  $(b \wedge (b \rightarrow c)) \rightarrow c$

B. 1.  $\neg(b \wedge \neg b)$

By De Morgan's law  
or  $\neg b \vee b = T$

$$\Rightarrow \neg(F)$$

$$\Rightarrow T$$

This statement is tautology  
because every interpretation of 1 will give true

2.  $(b \wedge (b \rightarrow c)) \rightarrow c$

$$(b \wedge (\neg b \vee c)) \rightarrow c$$

$$((b \wedge \neg b) \vee (b \wedge c)) \rightarrow c$$

$$(F \vee (b \wedge c)) \rightarrow c$$

$$(b \wedge c) \rightarrow c$$

$$\neg(b \wedge c) \vee c$$

$$(\neg b \vee \neg c) \vee c$$

$$\Rightarrow \neg b \vee (\neg c \vee c) \rightarrow c$$

$$\Rightarrow \neg b \vee T$$

$$\Rightarrow T \text{ (Tautology)}$$

$$\Rightarrow T \vee c$$

Q. Prove that

$$(\neg(b \rightarrow c) \wedge \neg(\neg b \rightarrow (c \vee d))) \rightarrow (\neg c \rightarrow d)$$

$$\Rightarrow (\neg(\neg b \vee c) \wedge \neg(\neg b \rightarrow (c \vee d))) \rightarrow (\neg c \rightarrow d)$$

$$(\neg(\neg b \vee c) \wedge \neg(b \vee (c \vee d))) \rightarrow (c \vee d)$$

$$(b \wedge \neg c) \wedge (\neg b \wedge \neg(c \vee d)) \rightarrow (c \vee d)$$

$$(b \wedge \neg c) \wedge (\neg b \wedge (\neg c \wedge \neg d)) \rightarrow (c \vee d)$$

$$(\underline{b} \wedge \neg c \wedge \neg b \wedge \neg c \wedge d) \rightarrow (c \vee d)$$

$$F \rightarrow (c \vee d)$$

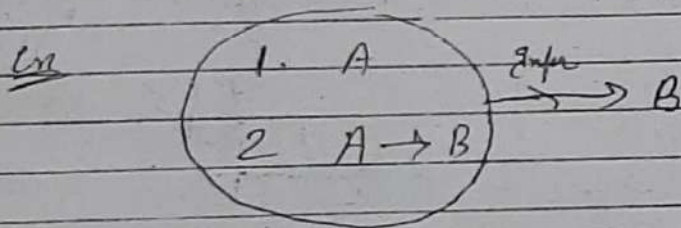
$$\neg(F) \vee (c \vee d)$$

$$\neg \neg \vee (c \vee d)$$

$$\underline{T} \quad (\text{Tautology})$$

### Reasoning Techniques (Inferencing Rules):

#### 1. Modus Ponens



#### 2. Principle of resolution: → - this

Inferencing technique is applicable if the given formulas are in clause form.

Clause form → A formula is said to be in clause form if only 2 operators are used i.e. '∨' and '¬'.  
 ¬ must be associated with individual proposition (not in group)

ex 1. A ∨ ¬B is clause  
 2. ~~but~~ ¬(A ∨ B) is not clause then we

break it. But  $\neg(A \vee B)$   
 $\Rightarrow \neg A \wedge \neg B$   
 $\Rightarrow \neg A$  & then  
 $\Rightarrow \neg B$  } These are clause.

#### 3. ¬(A ∨ B)

PqR.

ex 1. ¬A ∨ B } 3. B  
 2. A }

ex 2 1. A ∨ B } → 3. B  
 2. ¬A }

ex 3 1. A ∨ B } 3. B ∨ C  
 2. ¬A ∨ C }

#### 3. Chain Rule: →

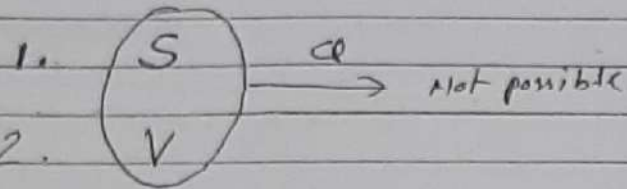
ex 1. A → B } 3. A → C  
 2. B → C }

Q 1. All ~~USE~~ SATI students are intelligent

2. A vijay is SATI student

Query  $\rightarrow$  Is vijay intelligent?

A



Query  $\rightarrow Q$

We can't do it by propositional logic.

Because propositional logic is very limited.

Predicate Logic (1<sup>st</sup> order logic)  $\rightarrow$

Predicate logic includes all the things which are the part of propositional logic formula.

In predicate logic there are extra facilities over propositional logic for knowledge representation these are  $\rightarrow$

i) Predicates can also be ~~imp operators~~

where predicate represents English sentences very similar to propositions but the structure of predicate is similar to the functions used in ~~just~~ high level languages.

ex.

Greater  $(x, y)$  :  $x$  is greater than  $y$

FATHER  $(x, y)$  :  $x$  is father of  $y$

FATHER  $(x)$  :  $x$  is father

Predicate will always be true or false.

(ii) Functions can also be used in predicate logic formula.

The structure of func<sup>n</sup> is very similar to structure of predicate but function returns non-boolean values.

Each func<sup>n</sup> represents English sentence.

ex.

Age  $(x)$  : age of 'x'

function

salary(x): Salary of 'x'

father(x): Father of 'x'

~~x~~  $x+y$  is greater than  $z$   
by predicate L.F.

~~e~~ GREATER (add(x,y), z)

add(x,y)  $\rightarrow$  x+y

(iii) In predicate logic formula  
quantifiers can also be used.

There are 2 types of quantifiers

a. Universal Quantifier  $\rightarrow \forall$  (For all)

b. Existential Quantifier

$\exists$  : there exist

$\exists$   $\forall$  STUDENT(x)

: Some are  $\forall$  students

~~Q~~ Convert the following sentences into  
predicate L.F.

(i) If  $x$  is on top of  $y$ ,  $y$   
supports  $x$ .

(ii) If  $x$  is above  $y$  and they are  
touching each other,  $x$  is on top of  $y$

(iii) A cup is above a book.

(iv) A cup is touching a book.

By resolution prove that book  
supports cup.

~~E~~ Following predicates are used

1. TOP(x,y) :  $x$  is on top of  $y$

2. SUPPORT(x,y) :  $x$  supports  $y$

3. ABOVE(x,y) :  $x$  is above  $y$

4. TOUCH(x,y) :  $x$  &  $y$  are touching  
each other.

(i) TOP(x,y)  $\rightarrow$  SUPPORT(y,x)

(ii) ABOVE(x,y)  $\wedge$  TOUCH(x,y)  $\rightarrow$  TOP(x,y)

(iii) ABOVE(cup, book)

(iv) TOUCH(cup, book)

~~By~~ Resolution  $\rightarrow$  we have to prove -  
book supports cup

$\forall$  SUPPORT(book, cup)

For <sup>propositional</sup> resolution we will have to make the predicates in clause form

Clauses:  $\rightarrow$

(i)  $\neg \text{TOP}(x, y) \vee \text{SUPPORT}(y, x)$

(ii)  $\neg (\text{ABOVE}(x, y) \wedge \text{TOUCH}(x, y)) \vee \text{TOP}(x, y)$

~~$\neg (\neg \text{ABOVE}(x, y) \vee \neg \text{TOUCH}(x, y)) \vee \text{TOP}(x, y)$~~

(iii)  $\text{ABOVE}(\text{cup}, \text{book})$

(iv)  $\text{TOUCH}(\text{cup}, \text{book})$

Apply resolution b/w clause (ii) & (iii) with  $x = \text{cup}$  &  $y = \text{book}$

Thus  $\rightarrow$  (v)  $\neg \text{TOUCH}(\text{cup}, \text{book}) \vee \text{TOP}(\text{cup}, \text{book})$

Apply resolution b/w clause (iv) & (v)

Now clause

(vi)  $\text{TOP}(\text{cup}, \text{book})$

Apply resolution b/w (i) & (vi) with  $x = \text{cup}$  &  $y = \text{book}$

(vii)  $\text{SUPPORT}(\text{book}, \text{cup})$

Q 1 Represent following sentences in predicate logic formula.

(i) All employees having salary more than 1 lac Rs. per annum supposed to pay income tax.

(ii) Some employees are sick today.

(iii) No employee earns more than the president.

Let Employee(x)  $\wedge$  President(y)

~~Predicates:  $\rightarrow$   $\neg \exists$  parameter (salary, sal(x))~~

~~1. SALARY(x) Employee have salary~~  
~~2. SALARY~~

(i)  $\forall x \text{ EMPLOYEE}(x) \wedge (\text{GREATER}(\text{sal}(x), 1L) \rightarrow \text{IT}(x))$

Following predicates are used:

$\text{EMPLOYEE}(x)$ : x is employee

$\text{GREATER}(x, y)$ : x is greater than y

$\text{IT}(x)$ : x pays income tax

Following functions are used:

$\text{salary}(x)$   $\text{sal}(x)$ : Salary of x

Q. (ii)  $\exists x \text{ EMPLOYEE}(x) \wedge \text{SICK}(x)$

Predicates  $\rightarrow$

$\text{SICK}(x)$ :  $x$  is sick

Q. 2 Make predicate L.F.

1. All SATI students are intelligent.
2. Amit is SATI student.

B

1. S (All students of SATI are intelligent)
2. A (Amit is SATI student)

Following predicates are used  $\rightarrow$

$\text{INTEL}(x)$ :  $x$  is intelligent

$\text{SATISTU}(x)$ :  $x$  is SATI student

1.  $\forall x (\text{SATISTU}(x) \rightarrow \text{INTEL}(x))$   
Here this is a rule so we use  $\rightarrow$  sign.
2.  $\text{SATISTU}(\text{Amit})$ .

Q. 3

Q. 1 (iii) No employee earn more than the president.

Predicates  $\rightarrow$

$\text{EMP}$   
 $\text{EMPLOYEE}(x)$ :  $x$  is employee

Function  $\rightarrow$   $\text{EARNING}(x)$ : earning of  $x$

$\text{PRES}(y)$ :  $y$  is the president

$\text{MORE}(x, y)$ :  $x$  is more than  $y$

$\forall x \forall y \text{ EMP}(x) \wedge \text{PRES}(y) \rightarrow \neg \text{MORE}(\text{EARNING}(x), \text{EARNING}(y))$

Q. 3

1. All the persons of age more than 60 years are Senior citizens

B Predicates  $\rightarrow$

$\text{PERSON}(x)$ :  $x$  is person

~~$\text{PERSON}(x)$~~

$\text{SCIT}(x)$ :  $x$  is Senior citizen

$\text{MORE}(x, y)$ :  $x$  is more than  $y$

Functions  $\rightarrow$

$\text{AGE}(x)$ : age of  $x$

~~$\text{AGE}(x)$~~

$\forall x (\text{PERSON}(x) \wedge \text{AGE}(\text{AGE}(x), 60) \rightarrow \text{SCIT}(x))$

Q. 4

Students get 60% or more in 12<sup>th</sup> and 35% marks in PET are eligible for engg. Admission

B

Predicates:

$\text{STU}(x)$ :  $x$  is student

$\text{MARK}(x)$ :  $M$

$\text{ELIGIBLE}$ :  $x$  is eligible for engg. Admission

$\text{MORE}(x, y)$ :  $x$  is more than  $y$



Functions:

$m$  marks<sup>-12</sup>( $x$ ): marks of  $x$  in 12<sup>th</sup>

Marks-PET( $x$ ): marks of  $x$  in PET

$\forall x$  STU( $x$ )  $\wedge$  MORE (marks-12( $x$ ), 60%)  
 $\wedge$  MORE (marks-PET( $x$ ), 35%)  
 $\rightarrow$  ELIGIBLE ( $x$ )

Conversion of predicate logic formula  
into Prece Normal Form (PNF):  $\rightarrow$

$(Q_1x) (Q_2x) \dots (Q_nx) M$

$Q_1, Q_2, \dots, Q_n$  are the quantifiers.

$M$  - It is the formula (P.L.F) that  
doesn't contain any quantifier.

Rules required to convert predicate logic  
formula into PNF:  $\rightarrow$

1.  $(Qx) F(x) \vee G = (Qx) [F(x) \vee G]$

$F(x)$  &  $G$  are the predicates and  
propositions respectively

$Q$  - Quantifier

ex  
then  $\frac{F(y) \vee (Q(x)) G(x)}{\Rightarrow (Qx) [F(y) \vee G(x)]}$

2.  $(Qx) F(x) \wedge G = (Qx) [F(x) \wedge G]$

3.  $\neg (\forall x) P(x) = (\exists x) \neg P(x)$

4.  $\neg (\exists x) P(x) = (\forall x) \neg P(x)$

5.  $\forall x F(x) \wedge \forall x H(x) = \forall x [F(x) \wedge H(x)]$

6.  $\forall x F(x) \vee \forall x H(x) \neq \forall x [F(x) \vee H(x)]$

7.  $\exists x F(x) \vee \exists x H(x) = (\exists x) [F(x) \vee H(x)]$

8.  $\exists x F(x) \wedge \exists x H(x) \neq \exists x [F(x) \wedge H(x)]$

9.  $\forall x F(x) \vee \forall z H(z) = \forall x \forall z [F(x) \vee H(z)]$

10.  $\exists x F(x) \wedge \exists x H(x) = \exists x \exists y [F(x) \wedge H(y)]$

11.  $(Q_1x) F(x) \vee (Q_2x) H(x) = (Q_1x) (Q_2y) [F(x) \vee H(y)]$

12.  $(Q_1x) F(x) \wedge (Q_2x) H(x) = (Q_1x) (Q_2y) [F(x) \wedge H(y)]$

convert it into PNF

$$\forall x \forall y [(\exists z) (p(x,z) \wedge p(y,z))$$

$$\rightarrow (\exists u) q(x,y,u)]$$

1. 2.

$$\Rightarrow (\forall x)(\forall y) [\neg (\exists z) (p(x,z) \wedge p(y,z)) \vee (\exists u) q(x,y,u)]$$

$$\Rightarrow (\forall x)(\forall y) [(\forall z) (\neg p(x,z) \vee \neg p(y,z)) \vee (\exists u) q(x,y,u)]$$

$$\Rightarrow (\forall x)(\forall y)(\forall z) (\neg p(x,z) \vee \neg p(y,z) \vee q(x,y,u))$$

(By 1<sup>st</sup> rule)

$$\forall x p(x) \rightarrow (\exists x) q(x)$$

$$\neg (\forall x p(x)) \vee (\exists x) q(x)$$

$$\Rightarrow \exists x \neg p(x) \vee (\exists x) q(x)$$

$$\Rightarrow (\exists x) [\neg p(x) \vee q(x)]$$

Skolem function:  $\rightarrow$

(Skolemization)

All  $\exists$  quantifiers should  
b. left most then followed by  
 $\forall$

ex:

$$\forall x \exists y [\neg p(x) \vee q(y)] = \exists y \forall x$$

$$[\neg p(x) \vee q(y)]$$

$$\exists x (\exists y) (\forall z) (\exists u) \forall (v) (\exists w) p(x,y,z,u,v,w)$$

skolemize it!

$$\exists x (\exists u) (\exists w) (\forall y) (\forall z) (\forall v)$$

$$p(x,y,z, f(y,z), v, g(y,z,v))$$

Conjunctive Normal Form (Group step to  
convert sentence into PNF):  $\rightarrow$

AND of OR

For ex.

1. A
2. A  $\vee$  B
3. A  $\wedge$  B
4. (A  $\vee$  B)  $\wedge$  ( $\neg$  C  $\vee$  D)

Disjunctive Normal form (OR of AND)

1. A
2. A  $\vee$  B
3. A  $\wedge$  B
4. (A  $\wedge$  B)  $\vee$  C

convert into CNF:  $\rightarrow$

$$1. (p \rightarrow q) \rightarrow r$$

$$(\neg p \vee q) \rightarrow r$$

$$\neg(\neg P \vee Q) \vee R$$

$$(P \wedge \neg Q) \vee R$$

$$\Rightarrow (\neg P \wedge R) \vee (P \vee R) \wedge (\neg Q \vee R)$$

2.  $P \vee (\neg P \wedge Q \wedge R)$

$$= (P \vee \neg P) \wedge (P \vee Q) \wedge (P \vee R)$$

$$= T \wedge (P \vee Q) \wedge (P \vee R)$$

$$= (P \vee Q) \wedge (P \vee R)$$

Conversion of Predicate logic formula into clause form:

1. Eliminate all implication & equivalence by law of implication & law of equivalence  
 $A \rightarrow B \Rightarrow \neg A \vee B$   
 $A \leftrightarrow B \Rightarrow (A \rightarrow B) \wedge (B \rightarrow A)$

2. Move all negations to immediately precedes operands by DeMorgan's law.

3. Convert the formula into Prenex Normal form

4. Apply Skolem function to make all existential quantifiers in prefix of the formula. Delete all existential quantifiers.

5. Convert the formula into conjunctive normal form.
6. Eliminate all universal quantifiers & partition the formula into clauses from AND operator.

Convert into clause form:

$$\exists x \forall y [\neg \exists z P(f(x), y, z) \vee (\exists u Q(x, u) \wedge \exists v R(y, v))] ]$$

$$\Rightarrow \exists x \forall y [\neg \exists z \neg P(f(x), y, z) \vee \exists u \exists v (Q(x, u) \wedge R(y, v))] ]$$

$$\Rightarrow \exists x \forall y [\exists z \exists u \exists v (\neg P(f(x), y, z) \vee Q(x, u) \wedge R(y, v))] ]$$

$$\Rightarrow \exists x \forall y \exists z \exists u \exists v (\neg P(f(x), y, z) \vee Q(x, u) \wedge R(y, v)) ]$$

$$\Rightarrow \exists x \exists z \exists u \exists v (\neg P(f(x), y, z) \vee Q(x, u) \wedge R(y, v)) ]$$

$$\forall y (\neg P(f(x), y, z) \vee Q(x, u) \wedge R(y, v)) ]$$

$$\forall y ((\neg P(f(x), y, z) \vee Q(x, u)) \wedge (\neg P(f(x), y, z) \vee R(y, v))) ]$$

1.  $\neg P(f(x), y, z) \vee Q(x, u)$

2.  $\neg P(f(x), y, z) \vee R(y, v)$

- Ex 1 Consider the following two sentences
1. John likes all kind of food.
  2. Apples are food.

5. Bill eats peanuts & is still alive.

6. Sue eats ~~everything~~ Bill eats.

1. Convert into F.O. predicate logic.

2. Prove by backward chaining John likes peanuts.

3. Convert the formula into clause form

4. Prove by resolution - John likes peanuts

Following predicates are used:→

1. FOOD(x): x is food.

2. LIKE(x,y): x likes y

3. EAT(x,y): x eats y.

4. KILLED(x): x is killed.

I →  $\forall x \text{ FOOD}(x) \rightarrow \text{LIKE}(\text{John}, x)$

II →  $\text{FOOD}(\text{Apple})$

III →  $\text{FOOD}(\text{chicken})$

IV →  $\forall x \text{ FOOD}(x) \wedge \text{EAT}(y, x) \rightarrow \neg \text{KILLED}(y)$

$\forall x \forall y \text{ EAT}(x, y) \wedge \neg \text{KILLED}(y) \rightarrow \text{FOOD}(y)$

here we assume that  
not killed means still alive

i.e.  $\neg \text{KILLED}(x) \leftrightarrow \text{ALIVE}(x)$

V  $\text{EAT}(\text{Bill}, \text{peanuts}) \wedge \neg \text{KILLED}(\text{peanuts})$   
or  $\text{ALIVE}(\text{Bill})$

VI  $\forall x \text{ EAT}(\text{Bill}, x) \rightarrow \text{EAT}(\text{Sue}, x)$

Remaining at last page

Q2 The law says it is a crime for an American to sell weapons to a hostile nation. The country Mono, an enemy of America has some missiles and all of its missiles were sold to it by Kernal West, who is an American.

(i) Convert the sentences in above paragraph into predicate logic statements.

(ii) Use resolution to do that West is criminal.

$\text{HOSTILE}(x) \leftrightarrow \text{EN}(x)$

Following predicates are used:→

$\text{ENEMY}(x)$ : x is enemy country of America

$\text{SELLS}(x, y)$ : x sells weapons to y

$\text{AMERICAN}(x)$ : x is an American

$\text{CRIMINAL}(x)$ : x is criminal

$\text{HAS}(x, y)$ : x has y

(i)  $\forall x \forall y \text{ American}(x) \wedge \text{SELL}(x, y) \wedge \text{ENEMY}(y) \wedge \text{SELLS}(x, y) \rightarrow \text{CRIMINAL}(x)$

II  $\text{ENEMY}(\text{Mono}) \wedge \text{AMERICAN}(\text{West}) \wedge \text{SELLS}(\text{West}, \text{Mono}, \text{weapon})$   
 ~~$\text{AMERICAN}(\text{West}) \wedge \text{HAS}(\text{Mono}, \text{weapon})$~~

(ii) We have to prove CRIMINAL(WEST)  
Convert it into clause form

(a)  $\neg$  American(x)  $\vee$   $\neg$  Enemy(y)  $\vee$   $\neg$  sell(x,y)  
 $\vee$  Criminal(x)

(b) ENEMY(Nono)

(c) American(west)

(d) sell(west, nono)

(e) Has(Nono, weapon)

(f)  $\neg$  Enemy(y)  $\vee$  sell(west, weapon)

Apply resolution b/w clause ①  
& ③ where  $x = \text{west}$

(g)  $\neg$  Enemy(y)  $\vee$   $\neg$  sell(west, y)  $\vee$  ~~Enemy~~ Criminal(west)

Apply resolution b/w ② & ①  
with  $y = \text{Nono}$ .

(h)  $\neg$  sell(west, nono)  $\vee$  Criminal(west)

(i) Criminal(west)

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Ques 2 Assume the following facts and rules

(i) Steve only likes easy courses.

(ii) Science courses are hard.

(iii) All the courses in electronics don't are easy

Enemy Steve(Steve, x)  
like

N7

Use resolution to answer the query what course would Steve like.

Ans

Following predicates are used:  $\rightarrow$

(i) EASY(x): x likes easy courses

(ii) HARD(x): x is hard course  
 $\Leftrightarrow \neg$  EASY

(iii) FLEX(x): x is electronics course

(iv) LIKE(x, y): x likes y

~~$\forall$  x FLEX(x):~~

①  $\forall$  x EASY(x)  $\rightarrow$  LIKE(Steve, x)

②  $\neg$  EASY(Science)

③  $\forall$  x FLEX(x)  $\rightarrow$  EASY(x)

④ FLEX(EL-301)

Clause form:  $\rightarrow$

①  $\neg$  EASY(x)  $\vee$  LIKE(Steve, x)

②  $\neg$  EASY(Science)

③  $\neg$  FLEX(x)  $\vee$  EASY(x)

④ FLEX(EL-301)

quantifiers are coming at left

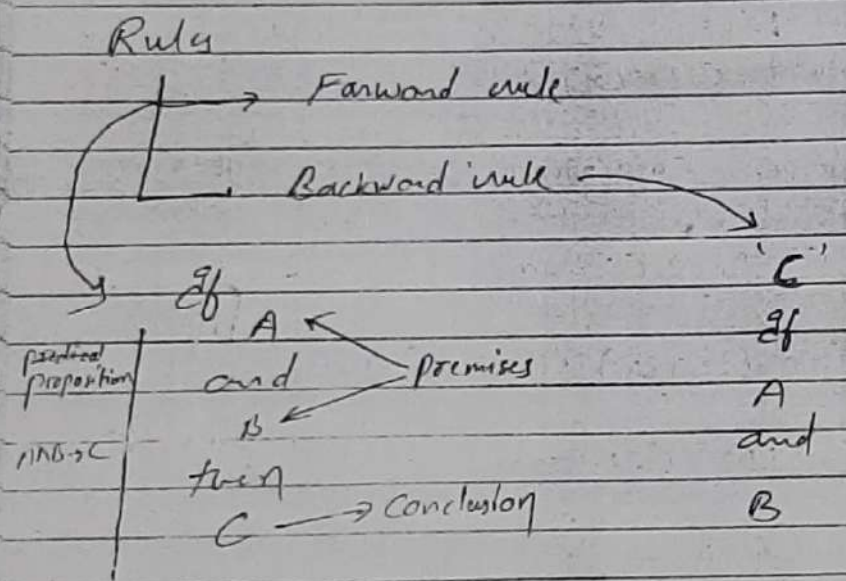
(5) EASY (EL-301)

Apply rule b/w (1) & (5)

(6) LIKE (Steve, EL301) (which was the query)

Backward Chaining →

Knowledge Represented by Rules & facts



ex 1.

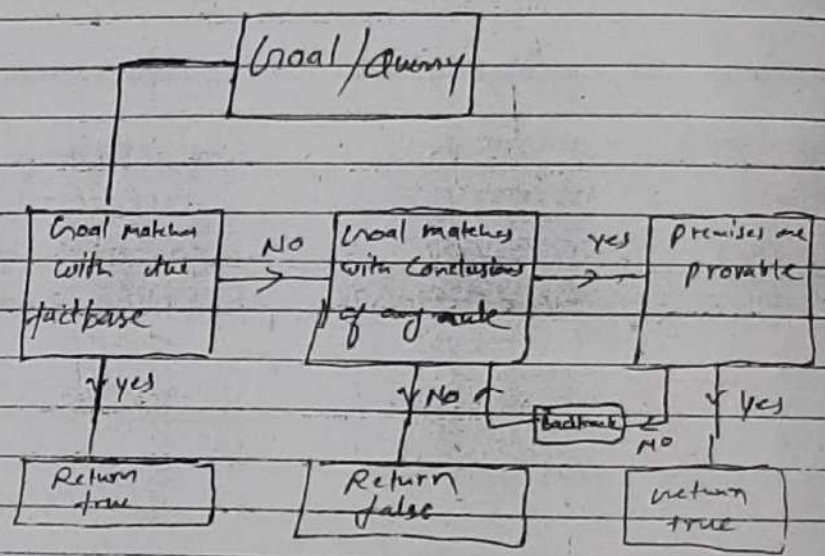
if Easy(x)

then Like (Steve, a)

Backward Reasoning / Chaining →

When the rules are written in the backward form and various rules and facts are changed for reasoning then this method is known as backward chaining.

Block Diagram of backward chaining →



ex Rules → (previous ex. rules)

(1) ~~Easy(x)~~ Like (Steve, a)

if Easy(x)

(2) Easy(x)

if ELEX(x)

facts: →

②  $\neg$  easy (science)

③ FLEX (FL-301)

Q: if query = Like (Steve, FL-301)

Query is matched with conclusion of rule-1

like (Steve, x)

↓  
easy(x) - {premise is not matched with any fact}

Try to match it with conclusion of other rules.

→ matched with conclusion of rule no-2

↓  
FLEX(x)

matched with fact no-2 with  $x = FL-301$

B1 Backwards chaining: →

1.  ~~$\forall x$~~  ~~LIKE~~ FOOD(x) → LIKE(John, x)

2. FOOD (Apple)

3. FOOD (chicken)

4.  ~~$\forall x \forall y$~~  EAT(x, y)  $\wedge$   $\neg$  KILLED(x) → FOOD(y)

5. EAT (Bill, peanuts)  $\wedge$   $\neg$  KILLED (Bill)

6.  $\forall x$  EAT (Bill, x) → EAT (Sue, x)

Facts: →

1. FOOD (Apple)

2. FOOD (chicken)

3. EAT (Bill, peanuts)  $\wedge$   $\neg$  KILLED (Bill)

Rules: →

1. FOOD(x) → LIKE (John, x)

2. EAT(x, y)  $\wedge$   $\neg$  KILLED(x) → FOOD(y)

3. EAT (Bill, x) → EAT (Sue, x)

Query is LIKE (John, peanuts)

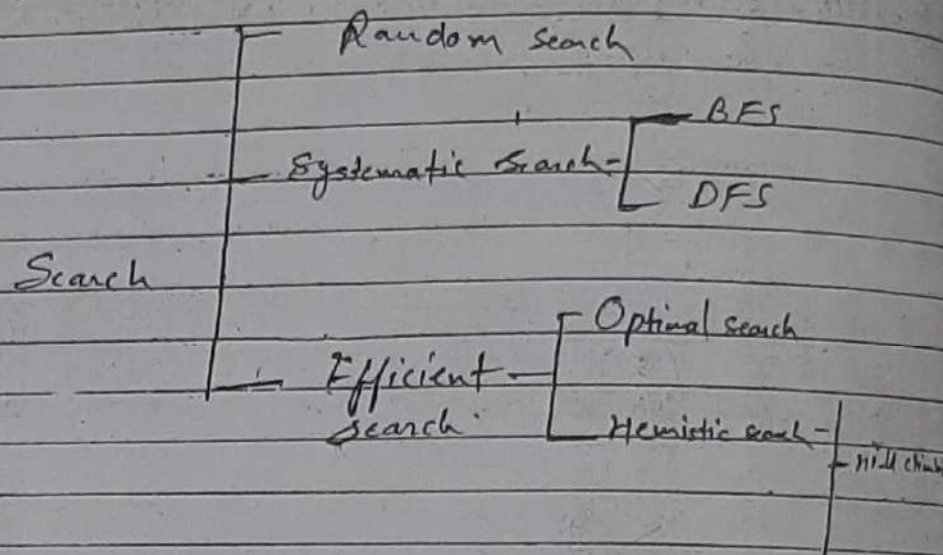
Rules 1. LIKE (John, x)

2. FOOD(y)  
 if  
 EAT(x,y)  
 and  
 ¬KILLED(x)

3. EAT(Sue, x)  
 if  
 EAT(Bill, x)

Query → LIKE(John, peanuts)  
 ↓  
 FOOD(x)  
 ↓  
 EAT(x, peanuts)  
 ∧ ¬KILLED(x)  
 ↓ which is a  
 fact so the given  
 query is true.

## Searching



Complexity of any greedy technique can not be more than complexity of sorting.

Heuristic search: → Searching for best but if not found then compromise with good solution. We will concentrate on <sup>time</sup> complexity more.

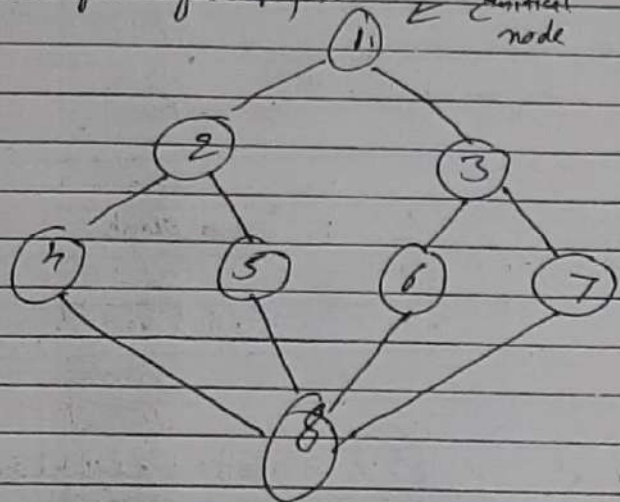
Traversal: → Searching of whole list means we will reach to last node.

Search: → We search till the required element is found.



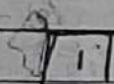
# Breadth First Search (BFS) $\Rightarrow$

State space dig (Graph)  $\Rightarrow$



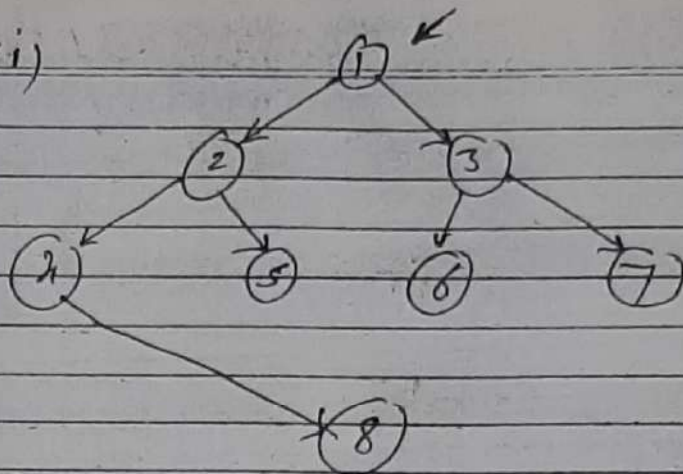
Here queue is used.

Here there is no left right nodes.

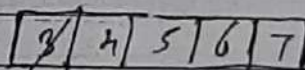
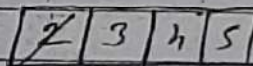
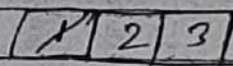


- Q. For the given graph if searching starts at node  $\textcircled{1}$  find
- (i) B.F. Spanning tree
  - (ii) Order of traversal.

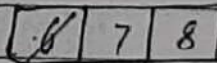
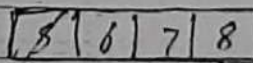
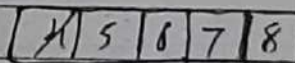
B (i)



Order: 1, 2, 3, 4, 5, 6, 7, 8

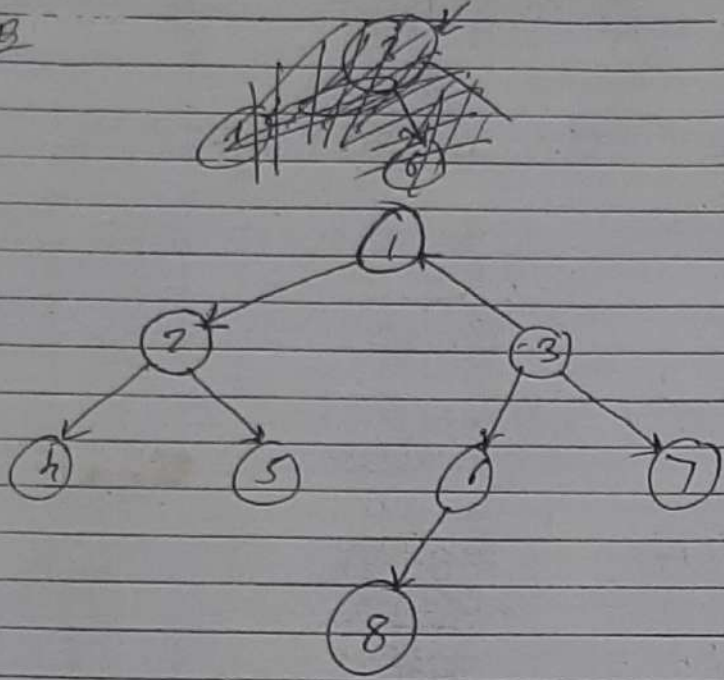


1 has been traversed



Q Start from node -3 and do same

B



Order: 3, 1, 6, 7, 2, 8, 4, 5

~~3~~ | 1 | 6 | 7

~~1~~ | 6 | 7 | 2 | 8

~~6~~ | 7 | 2 | 8

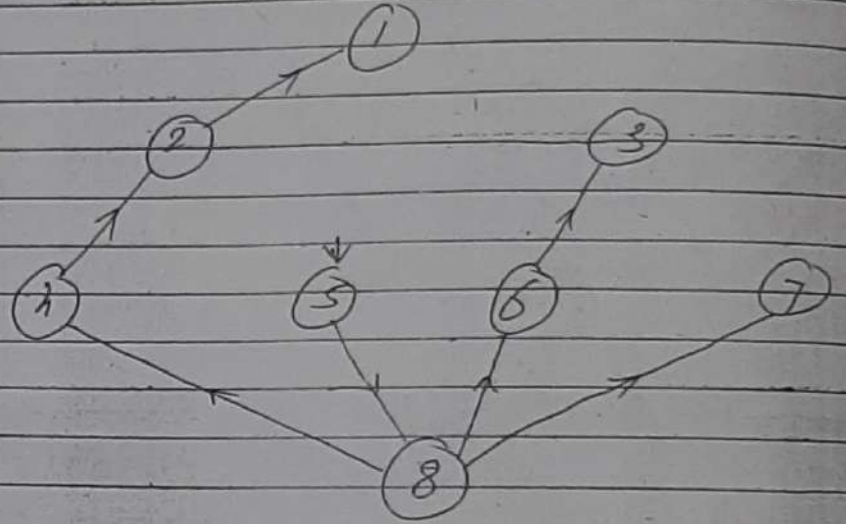
~~7~~ | 2 | 8

~~2~~ | 8 | 4 | 5

[ ]

Q Do same from (5) & (8)

Order: 5, 8, 4, 6, 7, 2, 3, 1



5 | | |

8 | | |

4 | 6 | 7

6 | 7 | 2

7 | 2 | 3 | 8

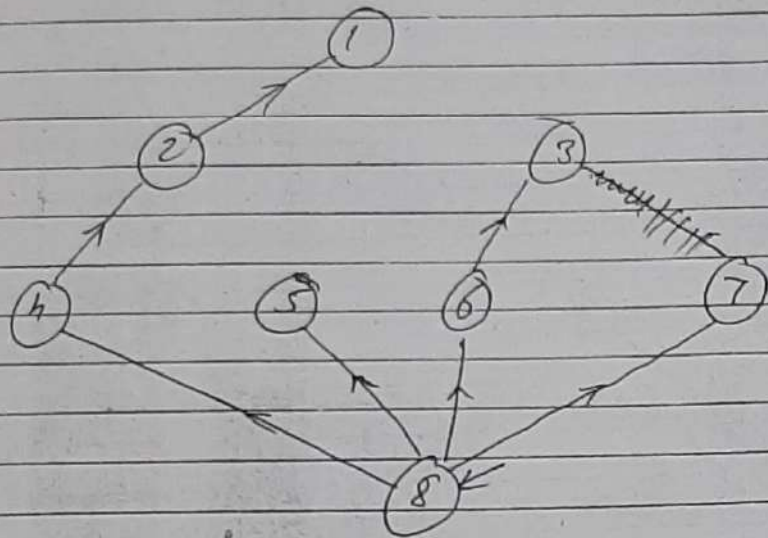
2 | 3 | |

3 | 1 | |

1 | | | → [ ]

From node 8 :->

Order :->



8			
---	--	--	--

4	5	6	7
---	---	---	---

5	6	7	2
---	---	---	---

6	7	2
---	---	---

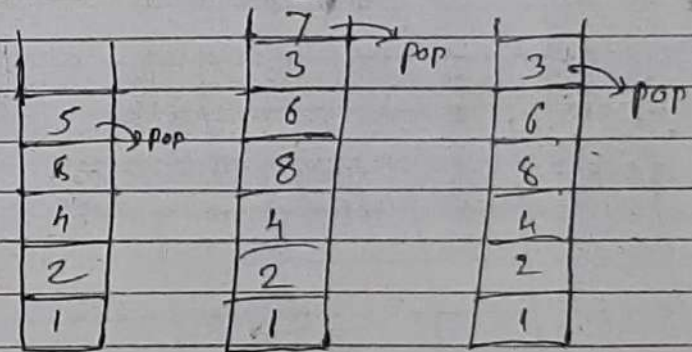
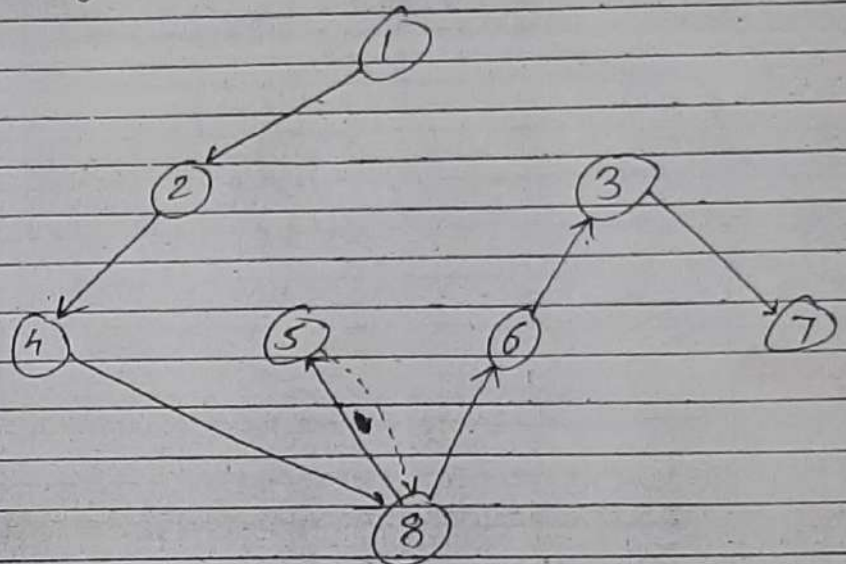
7	2	3
---	---	---

2	3	
---	---	--

3	1
---	---

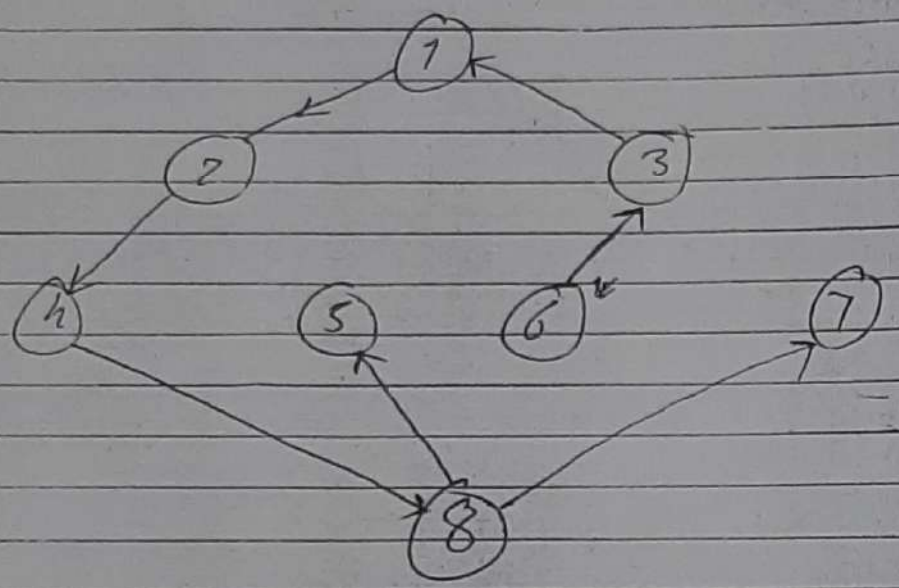
Depth first search (DFS) :->

Do ~~not~~ the same question by DFS  
Spanning tree :->



Order :-> 1, 2, 4, 8, 5, 6, 3, 7

Start from node 6



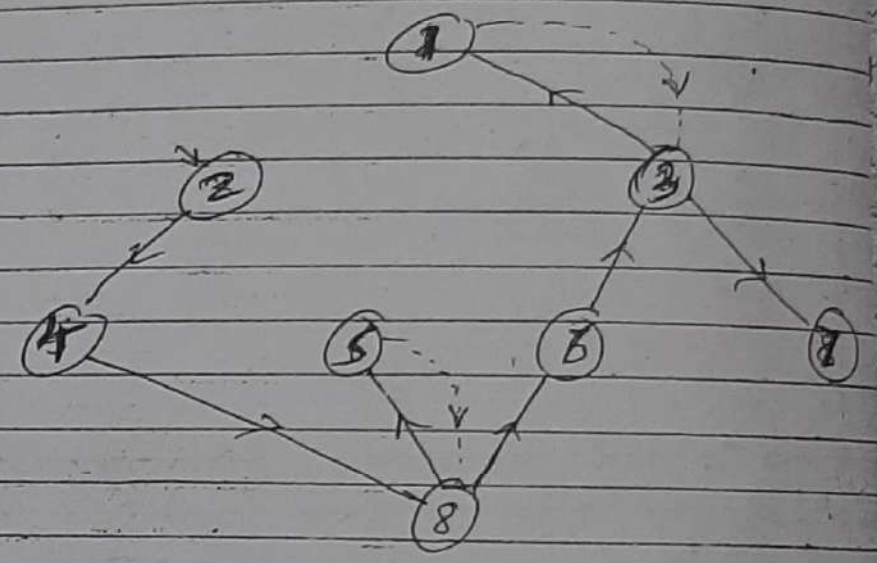
Order: 6, 3, 1, 2, 4, 8, 5, 7

	7		
5 → pop	8		
4			
2			
1			
3			
6			

Make it by 8 and 2

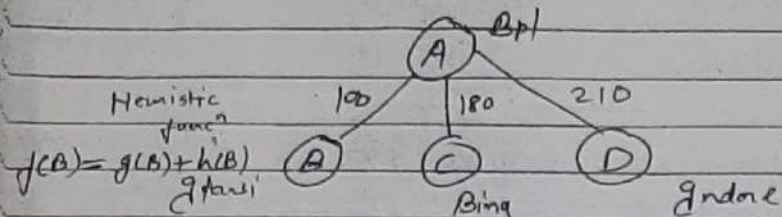
with 2

Order: 2, 4, 8, 5, 6, 3, 1, 7



	7				
1 → pop	3				
5 → pop	6			3	
8	8			6	6
4	4			8	8
2	2			4	4
				2	2

## Heuristic Search: →



(H) Goal (Chennai)

## Heuristic function: →

$$f(A) = g(A) + h(A)$$

where:

$f(A)$ : Total cost from node A to goal node

$g(A)$ : Actual cost to reach to node A from the initial node.

$h(A)$ : It is a heuristic cost of node A i.e. estimated cost calculated by applying heuristic function to reach the goal node.

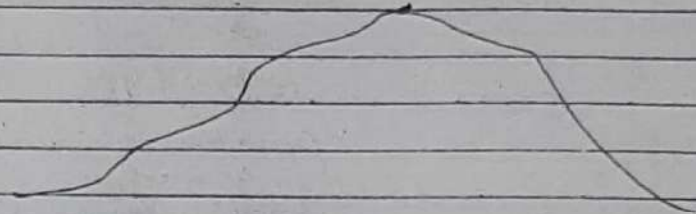
→ predicted cost

$$f(B) = 100 + 900 = 1000$$

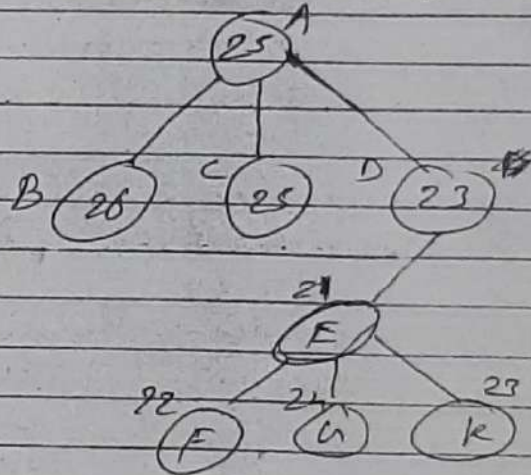
$$f(C) = 180 + 800 = 980$$

$$f(D) = 210 + 600 = 810$$

## Hill Climbing algo: →



Here we do the same as we do in case of making steps.



Here it will say "Sol" is cannot proceed.

## Algorithm: →

- Check whether initial node is the goal node if yes then stop with successful search otherwise calculate its total cost and make it the current node.

2 Repeat the following steps-

(a) Generate one possibility of current node.

(b) Check whether this node is goal if yes then stop with successful search.

(c) If no calculate the cost of the possibility (newly generated node).

If cost is better than the current then make new node as a current node and goto step (a)

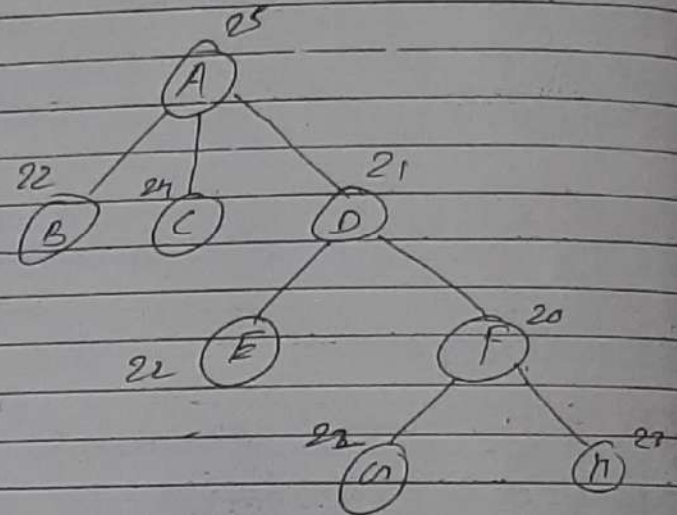
(d) Check whether the cost of newly generated node is worse than generate other option and do the same as of step (a) (b) (c)

If no option left then stop with unsuccessful search

Disadvantage

There is no guarantee of solution.

Steepest Hill Climbing:-



From algorithm point of view it is more promising algo.

So it will give better sol<sup>n</sup> than hill climbing.

Here time complexity will be more because we search for the best opt<sup>n</sup>

Algo →

(1) Check whether the initial node is the goal node if yes then stop with successful search.

Otherwise calculate the cost of and make it current node.

2. For the current node all possible options (children).

For each option do the following:  
 (a) If goal node is found stop with successful search.

(b) Calculate the cost of the nodes.

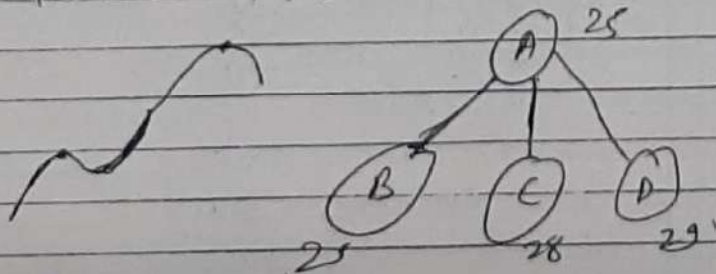
(c) Select best cost node (min cost) and compare its cost with current node.

If the cost is better than the ~~than~~ current node ~~then~~ make it current node and goto step-2.

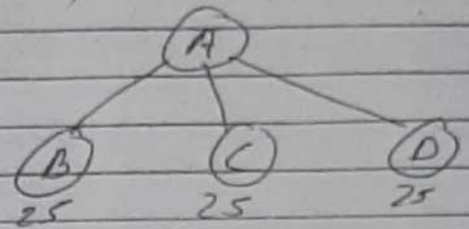
If cost is worse than the current then stop with unsuccessful search.

Problems & their sol<sup>n</sup> associated with these algos:

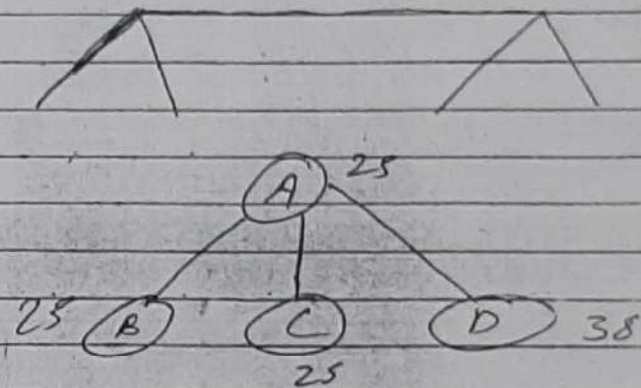
① Local maxima: →



② plateau  
 Plateau:



③ Ridge:-



Solutions: →

① Backtracking.

② Ignore some steps

But with these sol<sup>n</sup>s also there is no guarantee of finding the goal.

And this may also lead very high complexity ( $O(2^n)$ ).

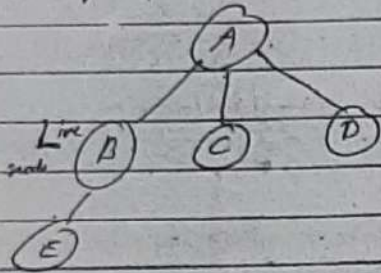
Branch and bound  $\rightarrow$

One of the examples of it is Best First Search.

Best First Search  $\rightarrow$

Live Nodes (Open list)  $\rightarrow$

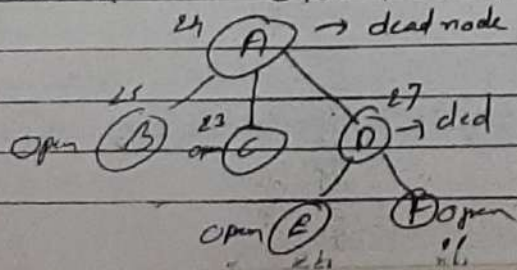
Nodes which are either partially expanded or not expanded at all are known as live nodes. Set of live nodes (list) is known as open list.



Dead nodes (closed list)  $\rightarrow$

Nodes which are either fully expanded <sup>and</sup> no further expansion is possible are known as dead nodes.

Set of dead node (list) is known as closed list.



so

closed list = {A, D}

Open list = {B, C, E, F}

E-Node  $\rightarrow$

Node selected from the open list which is most promising for further expansion is known as E-node.

E-Node = {C}

Algorithm  $\rightarrow$

Step 1 - Check whether initial node is goal node if yes then stop with successful search. Otherwise calculate cost and place it in open list.

2. Select most promising node from open list (least cost node). Selected node becomes E-node.

do the following.

a. Generate all possibilities from E-node and place it in closed list and it becomes dead node.

Do the following for each newly generated node

(a) check whether the generated node is goal node then stop with successful search.

(b) Otherwise calculate the cost of the node.



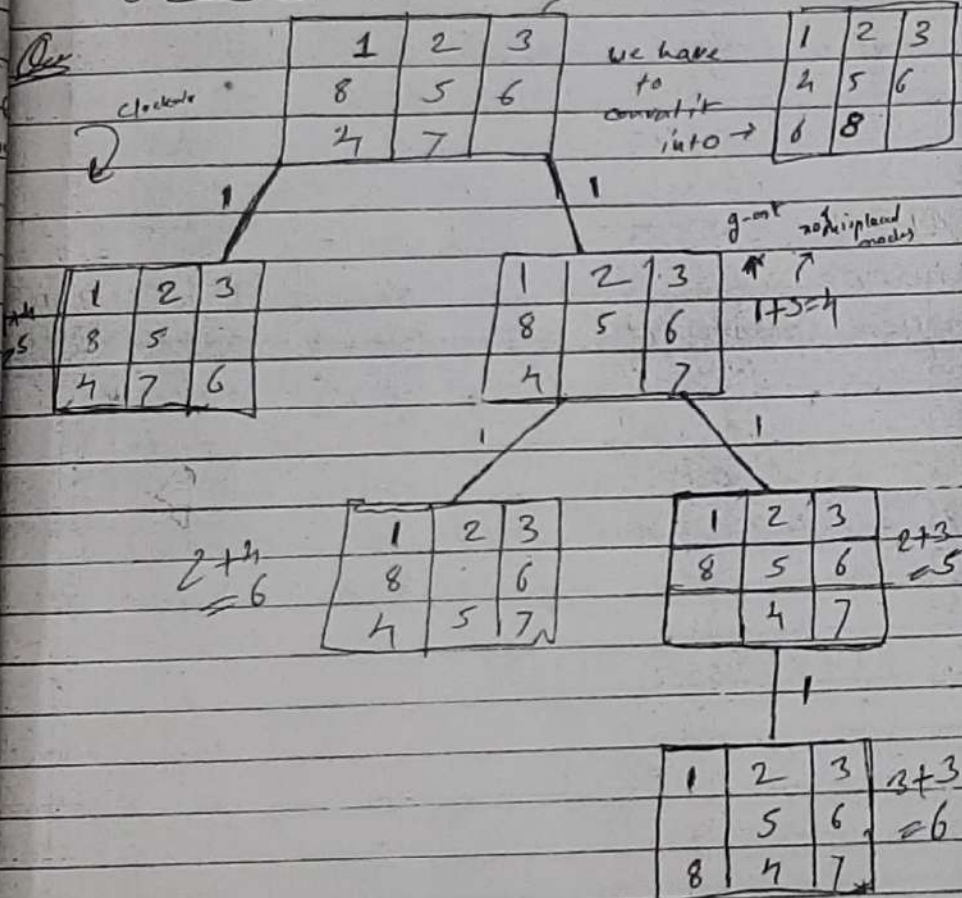
and place it in open list.

(c) goto step - 2

3. If open list is empty then stop with unsuccessful search.

8-puzzle problem:-

cost = 0 + 3 = 3



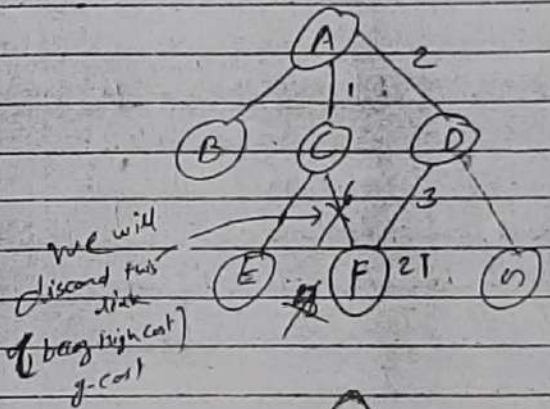
In previous example there was one initial & 1 goal?

- ① 1-initial, 1-goal → Ludo, Snake & ladders
- ② Many - ~~many~~ <sup>one</sup> → 15 puzzle, 24 puzzle
- ③ ~~Many~~ 1 to many → Chess, S.
- ④ Many-many → Tic-tac-toe, sudoku

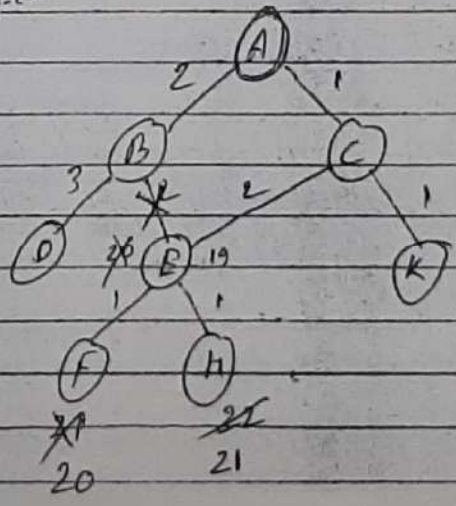
In 1-1 you can proceed from any one

A\* Algorithm:-

State space diagram is a graph here



here the children is already in open list



here the child is already in close list

Here we will do same as upper above case and

a node if it exists in state space diag

A\*

Algorithm →

It is one of the branch and bound techniques and very similar to best first search the difference is <sup>in BFS</sup> the state space diag. is around as tree but in A\* algo state space diagram may be graph. So A\* algo is more general version of Best First Search which is applicable for not only tree but also for graph.

Algo: →

(1) Check whether the start node or initial node is goal node if yes then stop with successful search otherwise calculate its cost and make it current node place it into open list.

(2) Generate select best option (least cost node) from open list generate all its children and for each child do the following →  
(If open list is empty then stop with unsuccessful search)

(a) Check whether the node (child) is goal node then stop with successful search else

(b) Calculate its cost

(c) Is the node already in open list, if yes then compare its new cost with already stored cost.

If new cost is better then change the parent of the node as per new path and delete old cost and restore new cost see the previous diag.

(d) Check whether the node already in closed list then compare its new cost with old cost, if the new cost is less

(i) Change the parent link to new path

(ii) Restore new cost and delete old cost.

(iii) Change the cost of all descendants see diag (2nd to show)

(e) Place the node in open list and goto step no. 2.

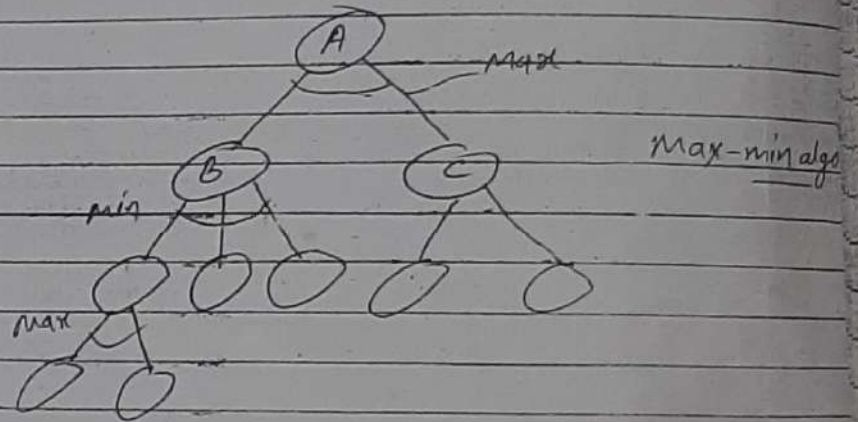
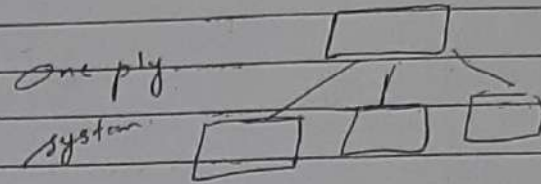
$$P(H|E) = \frac{P(H \cap E)}{P(E)}$$

$$P(E|H) = \frac{P(E \cap H)}{P(H)}$$

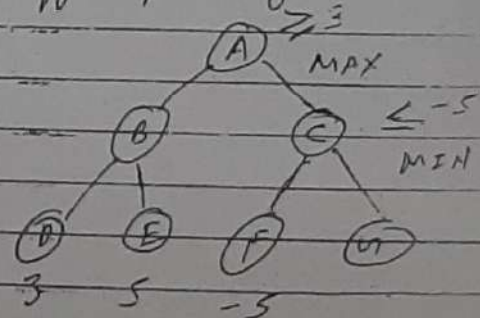
# Game Playing

## Max-min Technique

→  $\alpha$ - $\beta$  cutoff (Pruning)

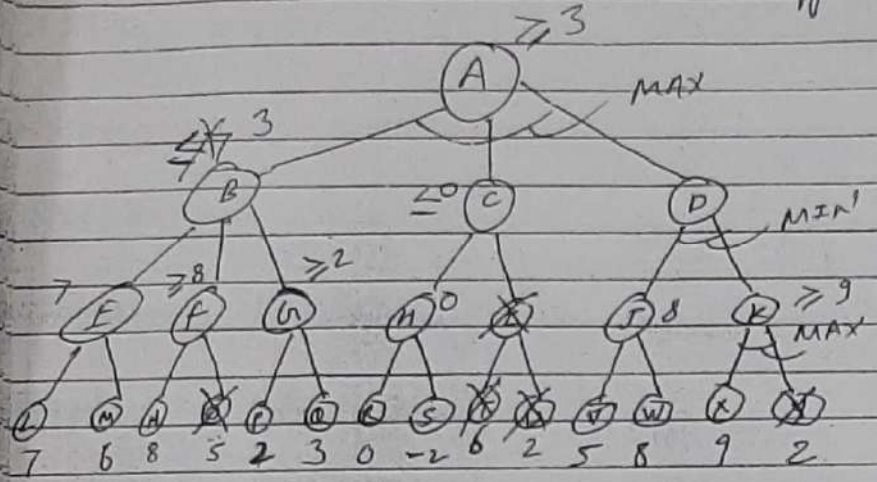


$\alpha$ - $\beta$  cutoff (Pruning) →

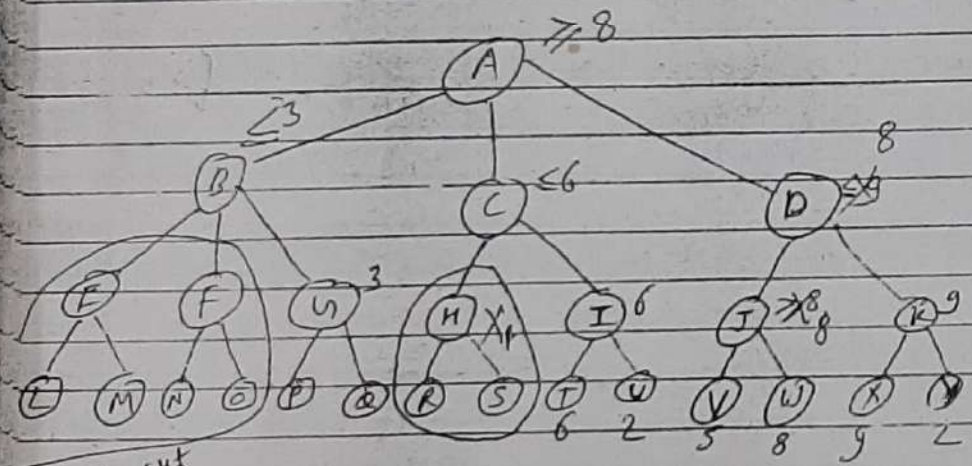


lower bound -  $\alpha$   
 upper bound -  $\beta$

Calculate  $M$  by max-MIN &  $\alpha$ - $\beta$  cutoff =



Left to Right cutoff.  
 so cutted nodes are  $I, J, K, Y$



Right to left cutoff

Unification Algorithm →

- Clauses:-
- $P(x, a) \vee Q(g)$
  - $\neg P(b, g)$
- Apply resolution b/w cl. 1 & 2  
 with  $x=b$  &  $a=g$
- $Q(a)$

This process is called term unification.

Algo

Input  $W = \{P(a, x), P(a, b)\}$

Output → 'σ' = {a/x, b/y}  
 (most general unifier) (mgu)  
 (mgu)

Algorithm :-

- Step-1 Set  $k=0, W_k=W$  &  $\sigma_k = \phi$
- Step-2 If  $W_k$  is singleton (A set whose cardinality is 1), stop,  $\sigma_k$  is mgu of  $W$ .  
 else find the Disagreement Set  $D_k$
- Step-3 If there exist elements  $V_k = d$  in  $D_k$  such that  $V_k$  is variable that does not occur in  $t_k$ , goto step 4.  
 else stop,  $W$  is not unifiable.
- Step-4  $\sigma_{k+1} = \sigma_k \cup \{t_k / V_k\}$

5.  $K = K + 1$ , goto step 2.

ex.  $W = \{P(a, x, f(g(y))), P(z, d(z), f(u))\}$

1.  $K = 0, W_0 = W, \sigma_0 = \{ \}$
2.  $D_0 = \{ a, z \}$
3.  $V_0 = z, t_0 = a$
4.  $\sigma_1 = \{ \} \cup \{ a/z \} = \{ a/z \}$   
 $W_1 = \{ P(a, z, f(g(y))), P(a, d(a), f(u)) \}$
5.  $K = 1$

part 2

2.  $D_1 = \{ x, f(a) \}$
3.  $V_1 = x, t_1 = f(a)$
4.  $\sigma_2 = \{ a/z, f(a)/x \}$
5.  $W_2 = \{ P(a, f(a), f(g(y))), P(a, f(a), f(u)) \}$
5.  $K = 2$

part 3

2.  $D_2 = \{ f(g(y)), f(u) \}$
3.  $V_2 = u, t_2 = g(y)$
4.  $\sigma_3 = \{ a/z, f(a)/x \} \cup \{ g(y)/u \}$
5.  $W_3 = \{ P(a, f(a), f(g(y))), P(a, f(a), f(g(y))) \}$

~~$W = \{ \theta(f(a), g(x)), \theta(y, z) \}$~~

part 1

1.  $K = 0, W_0 = W, \sigma_0 = \{ \}$
2.  $D_0 = \{ f(a), y \}$
3.  $V_0 = y, t_0 = f(a)$
4.  $\sigma_1 = \{ f(a)/y \}$
5.  $W_1 = \{ \theta(f(a), g(x)), \theta(f(a), z) \}$
5.  $K = 1$

part 2

2.  $D_1 = \{ g(x), z \}$
3.  $V_1 = z, t_1 = g(x)$
4.  $\sigma_2 = \{ g(x)/z \}$
5.  $W_2 = \{ \theta(f(a), g(x)), \theta(f(a), g(x)) \}$
5.  $K = 2$

Here there is no  $V_k$  and  $t_k$  exists for agreement so it is not unifiable.

# Planning :-

Divide and Conquer:- In divide & conquer we divide the problem in subproblems which are the same as the original problem but smaller.

while planning is divides the problem into the sequence of sub problems.

## Goal Stack Planning :-

### ex. Block world Problem:-

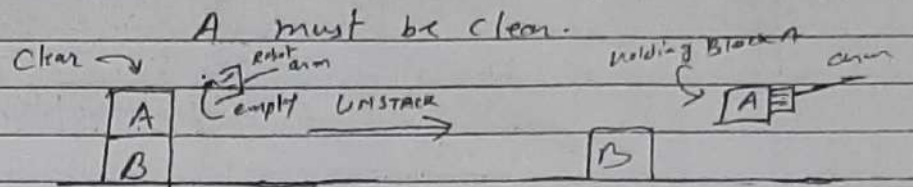
Predicates and proposition used to describe the state:-

1.  $ON(A, B)$ : Block A is on block B
2.  $ONTABLE(A)$ : A is on the table
3.  $CLEAR(A)$ : There is nothing on the top of A
4.  $HOLDING(A)$ : The Robot arm is holding block A
5.  $ARMEMPTY$ : Robot arm is holding nothing (empty)  
↳ Proposition

The actions performed by the Robot:-

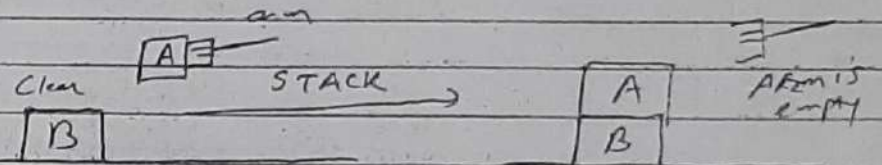
1.  $UNSTACK(A, B)$ : Pickup block A from its current position which is on B

Condition: Arm must be empty and block



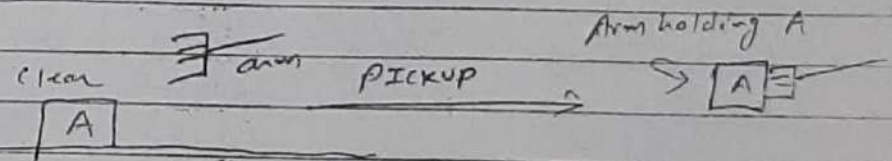
2.  $STACK(A, B)$ : Place block A on block B.

Cond: (i) The arm must already be holding block A. and  
(ii) Top of B should be clear



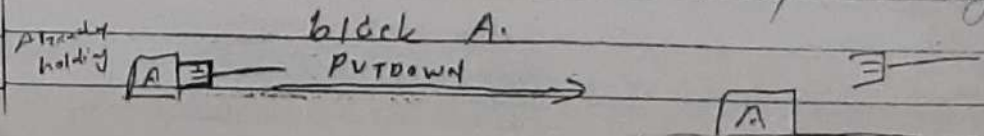
3.  $PICKUP(A)$ : Pick block A from the table and hold it.

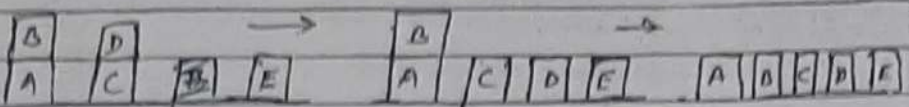
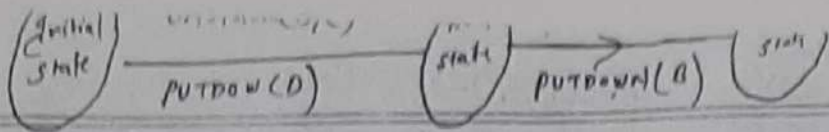
Cond: (i) The Robot arm must be empty.  
(ii) Top of A must be clear



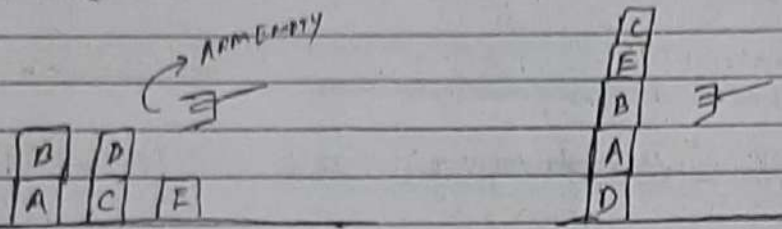
4.  $PUTDOWN(A)$ : Put block A down on the table.

Cond: Arm must be already holding block A.





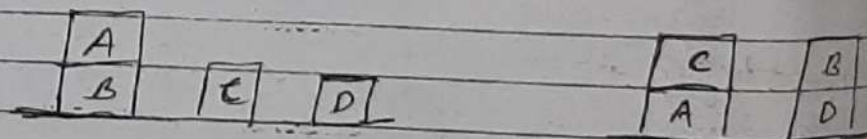
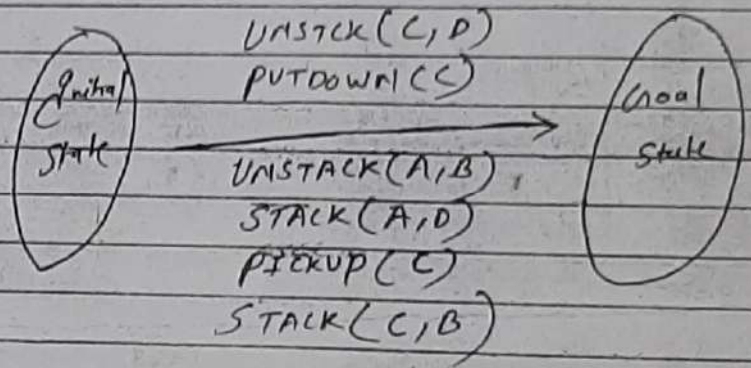
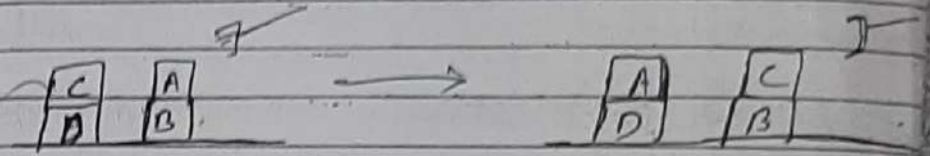
Initial state Goal state



- |  |   |  |
|--|---|--|
| ONTABLE(A)<br>ON(B,A) $\wedge$<br>ONTABLE(C) $\wedge$<br>ON(D,C) $\wedge$<br>ONTABLE(E) $\wedge$<br>CLEAR(B)<br>CLEAR(D)<br>CLEAR(E)<br>ARMEMPTY | UNSTACK(B,A)<br>PUTDOWN(B)<br>UNSTACK(D,C)<br>PUTDOWN(D)<br>STACKPICKUP(A)                  | ONTABLE(D)<br>ON(C,D) $\wedge$<br>ON(B,A) $\wedge$<br>ON(E,A) $\wedge$<br>ON(C,E) $\wedge$<br>CLEAR(C)<br>ARMEMPTY |
|  | STACK(A,D)<br>PICKUP(B)<br>STACK(B,A)<br>PICKUP(E)<br>STACK(E,B)<br>PICKUP(C)<br>STACK(C,E) |  |

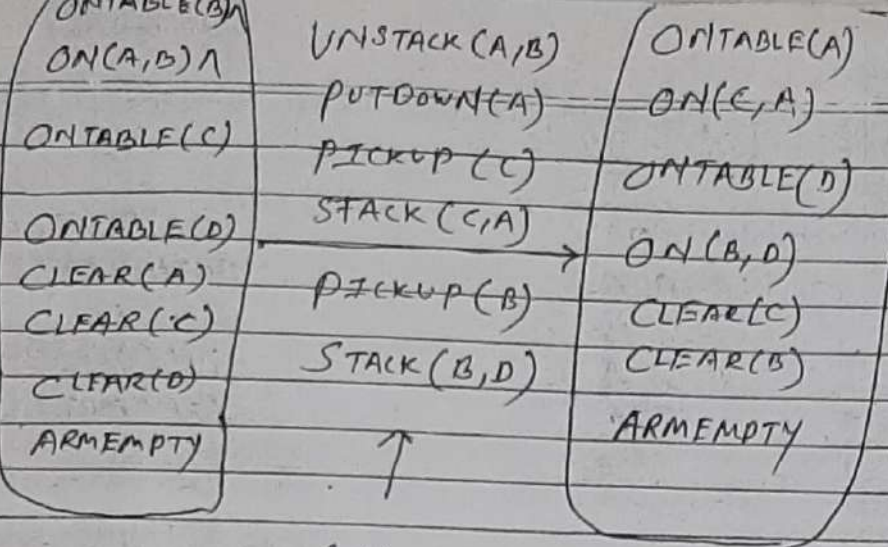
Initial state Goal state

- |  |  |
|--|--|
| ONTABLE(A) $\wedge$<br>ON(C,D) $\wedge$<br>ONTABLE(B) $\wedge$<br>ON(A,B) $\wedge$<br>ARMEMPTY | ONTABLE(D) $\wedge$<br>ONTABLE(C) $\wedge$<br>ON(A,D) $\wedge$<br>ON(C,B) $\wedge$<br>ARMEMPTY |
|--|--|



Initial state Goal state

To describe states predicates are used and for describing actions (transition) functions (actions) are used.



Initial                      Actions                      Goal

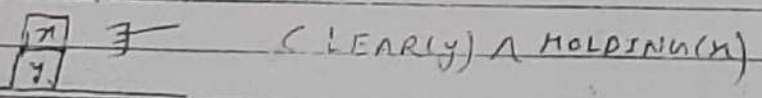
Q. Write precondition, Addlist, Deletelist for the actions considered in block world problem?

1. Preconditions  $\rightarrow$  (i) STACK(x,y)

b. ~~Stack~~: CLEAR(y)  $\wedge$  HOLDING(x)

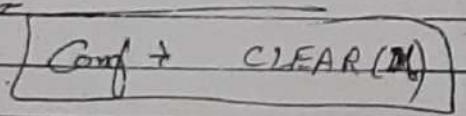


2. Deletelist  $\rightarrow$



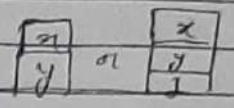
3. Addlist :-

ON(x,y)  $\wedge$  ARMEMPTY



(ii) UNSTACK(x,y)

Preconditions  $\rightarrow$



ARMEMPTY  $\wedge$  CLEAR(y)  $\wedge$  ON(x,y)

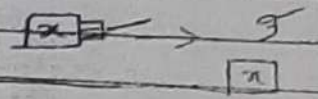
D:

ON(x,y)  $\wedge$  ARMEMPTY

A:

CLEAR(y)  $\wedge$  HOLDING(x)

(iii) PUTDOWN(x)



P:

HOLDING(x)

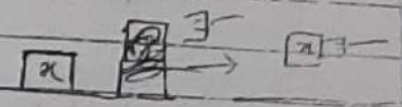
D:

HOLDING(x)

A:

ARMEMPTY  $\wedge$  ONTABLE(x)

(iv) PICKUP(x)  $\rightarrow$



P:

ARMEMPTY  $\wedge$  CLEAR(x)  $\wedge$  ONTABLE(x)

D:

HOLDING(x)  $\wedge$  ONTABLE(x)  $\wedge$  ARMEMPTY

A:

HOLDING(x)



## Knowledge Representation Techniques →

### 1. By Frame →

**FRAMES** → Frames are used to represent knowledge where declarations and procedures are combined into single knowledge representation environment. Frames are very similar to the objects, i.e. both data and procedures are packed into common structure.

**STRUCTURE OF FRAME** → A frame consists of:-

- (i) The "Name" of the frame → Each frame should have unique name.
- (ii) "Parent" of the frame → Frame uses concept of inheritance. The properties inherited by this frame is from its parent.
- (iii) "Slot names" and "their" "values" → Each frame may have large no. of slots. Within a frame the slot names must be distinct. But 2-different frames may have same "slot name" i.e. slot names are local to the frame. Each slot has some value associated with it.

Slots are considered to be similar as of fields of structure.

(iv) "Attached predicates" in each slot:- This is a unique property of the frame knowledge representation technique that each slot can be associated with 2- predicates. The purpose of these predicates are to provide security when the value of that slot is accessed or modified. These predicates are as follows →

- (a) "If Needed" predicate:- If a slot has an If Needed predicate then before the value of the slot is read and obtained, the If Needed predicate must be successfully proven (satisfied). The use of If Needed predicate is optional.
- (b) If added predicate:- If a slot has an If Added predicate, before the value of the slot is assigned or modified, its If Added predicate must be successfully proven. The use of this predicate is optional.

## Pictorial Representation of a Frame:

Name	Car
Parent	Automobile
Seat Seats	5
Engine	1335cc
Fuel	Diesel

we should write unique (specific) properties of car

Frame	Niloy
Parent	Teachers
Name	
Dob	
Salary	
Qual	

in Frames. (Feature of Frame)

By Script :- It is data representation technique which is used to represent seq<sup>n</sup>

Important components of the script are as follows

① Entry conditions :- Set of conditions that must be satisfied before the events stitely described in script occur.

② Result :- Conditions that will be true after the events described in script have occurred.

③ Props :- These are the physical objects which are essential parts of events described in script.

④ Roles :- People those are the part of the events/ scenes described in script are known as roles.

⑤ Track :- A same script may have different no. of tracks. These tracks share many but not all components of the script.

⑥ Scenes :- The actual seq<sup>n</sup> of events described in script.

known as scenes. Practically the scenes are represented using Conceptual Dependencies (CD).

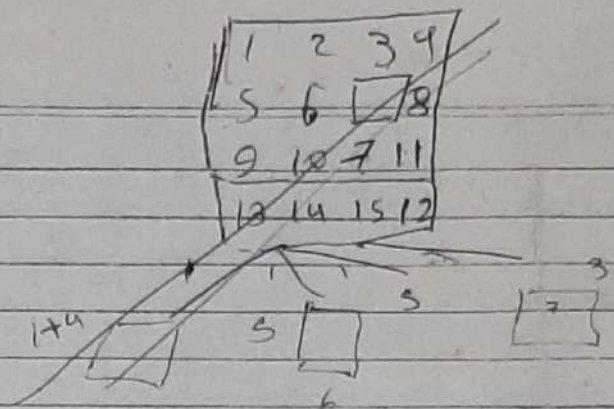
ex.

Script of Restaurant (Pictureially)

These are actually represented in CD but we write in English.

Script - Restaurant	Scene 1: Customer will go to rest.
Track: Coffee house	Scene 2: Customer will enter rest.
Props: Table, chair, food, money	Scene 3: Waiter will welcome him.
Roles: Manager, Customer, Manager, waiter, Cheefs.	Scene 4: Customer will sit to the table.
Entry conditions: Customer is hungry, Customer has money	Scene 5: Customer waiter will ask for services to provide.
Result: Customer is satisfied, Customer has less money, Manager will have more money	Customer will serve water, waiter will serve the required, Customer will eat.

our script university examination



Weak slot and filler  
Weak slot and filler

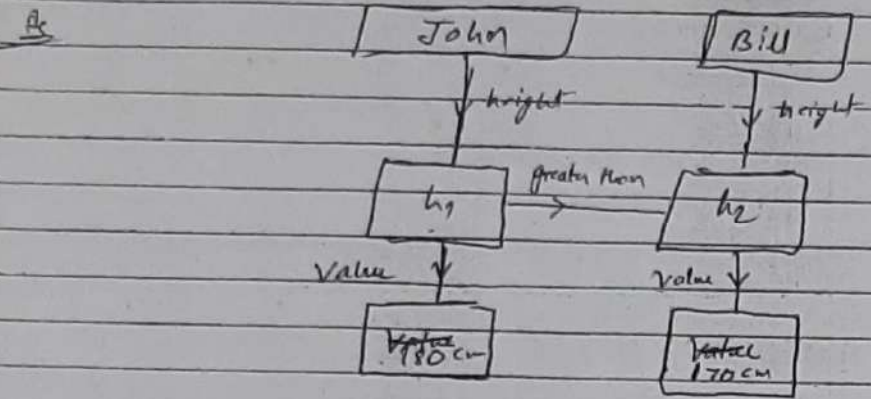
Semantic Net:-

Semantic net is a weak slot and filler knowledge representation technique. In weak slot and filler knowledge repr<sup>n</sup> different entities are connected to each other by infinite set of natural language terms. In this knowledge repr<sup>n</sup> it is very difficult to interpret correct meaning by diff. AT environments.

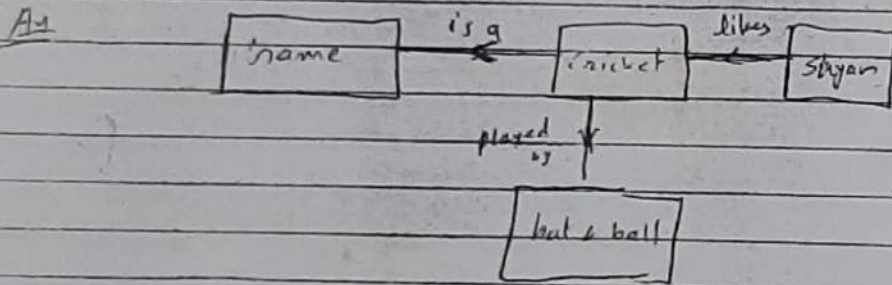
Q. Construct semantic net for following group of sentences

1. John's height is 180cm

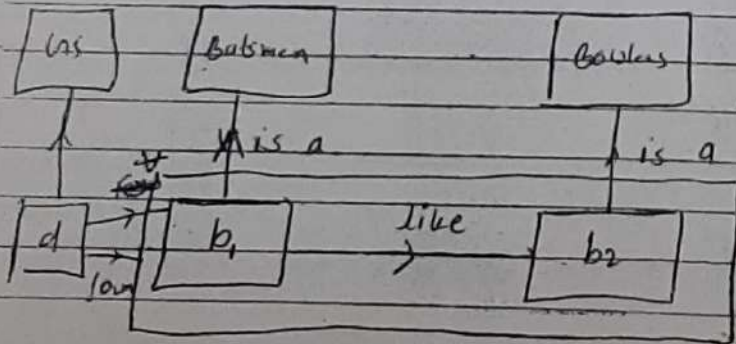
3. John is taller than Bill.



Q 1. Cricket is a game.  
2. It is played by bat and ball.  
3. Shyam likes cricket.

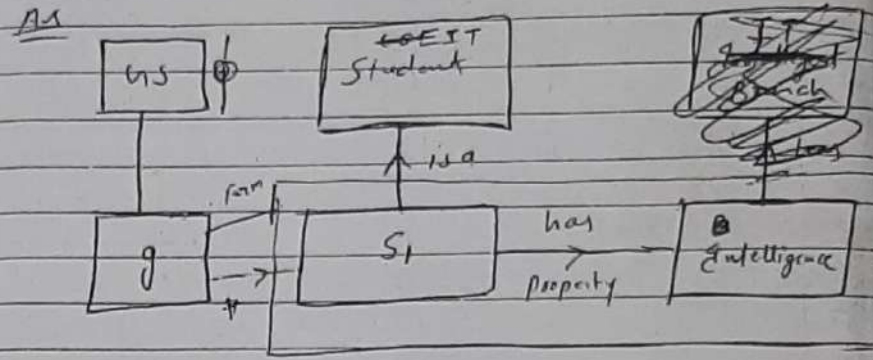


Q 1. All batsmen like batters, bowlers

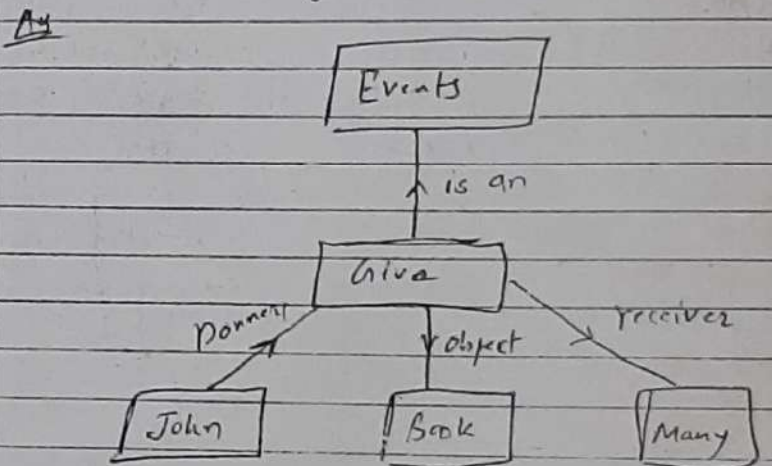


Q12. - General statements not related with other objects.

Q All IT Students are intelligent



Q John gave the book to many



Conceptual Dependency → It is one of the types of strong slot & filler structure.

C.D. → If the knowledge is represented in the form of semantic net i.e. weak slot & filler structure it is not universally acceptable i.e. the representation varies person to person.

The sol<sup>n</sup> for the above problem is to use strong slot & filler structure where activities used for rep<sup>n</sup> of knowledge are from finite & predefined set. The advantage of using conceptual dependency as knowledge rep<sup>n</sup> technique is:-

- ① Simpler methods to get inference stored knowledge base.
- ② The rep<sup>n</sup> is independent of the lang in which original knowledge is given, that's why this technique is universally acceptable technique.

The actions used in C.D. are as follows:-

ATRANS :- Transfer of an abstract relationship for example give and take.

PTRANS :- Transfer of physical location of an object.  
ex. go, come

PROPEL :- App<sup>n</sup> of physical force to the object for ex. push, pull etc.

MOVE :- Movement of body part by its owner. ex. kick

GRASP :- Grasping or holding an object by an actor.  
ex. Hold clutch

INGEST :- Ingestion of an object by any living thing.  
ex. Eat, Drink.

EXPEL :- Expulsion of something out from the body.  
ex. Cry

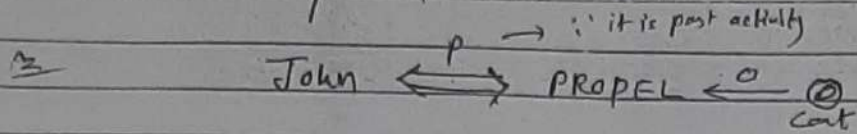
MTRANS :- Transfer of mental info.  
ex. Teaching, learning etc.

MBUILD :- Building new info. out of old ~~ex~~ inference or decide

SPEAK :- Production of sound

ATTEND :- Focusing of sense organ towards particular activity.  
ex - Listen, view

Q. John pushed the cart.

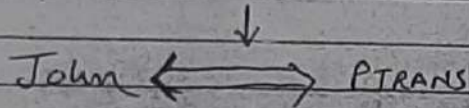


F - for future activity  
but noting for present

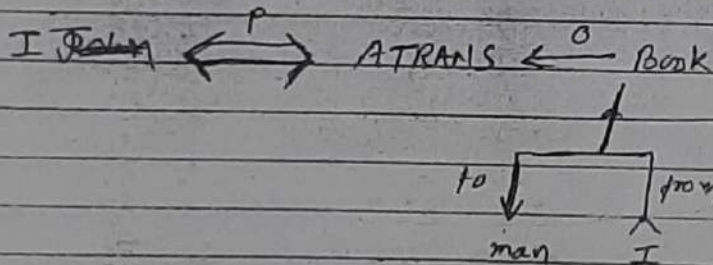
Q. 1. I gave book to a man

2. John come tomorrow.

A. 2. Tomorrow

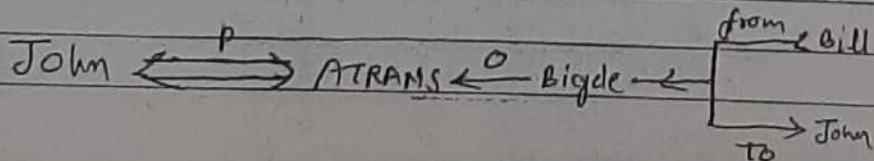


1.



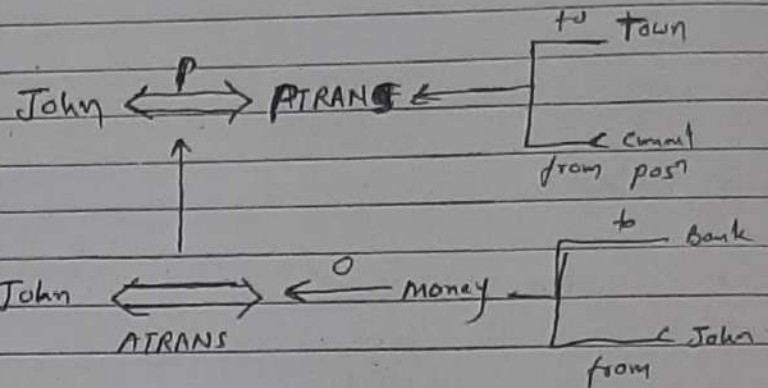
Q. John took a bike from Bill

A.



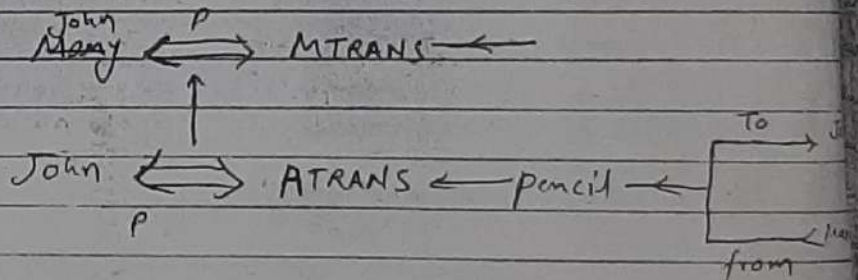
Q. John went down town to deposit money in the Bank.

A.



Q. John begged Mary for a pencil.

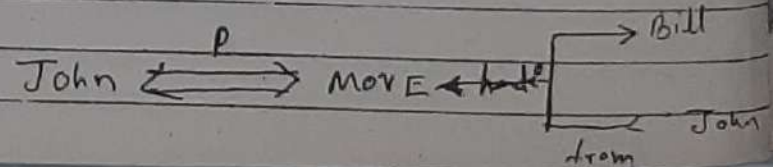
A.



Q. 1. John Slapped Bill

2. Bill drank cake.

A. 1.



John  $\leftrightarrow$  PROPEL  $\leftarrow^0$  Bill

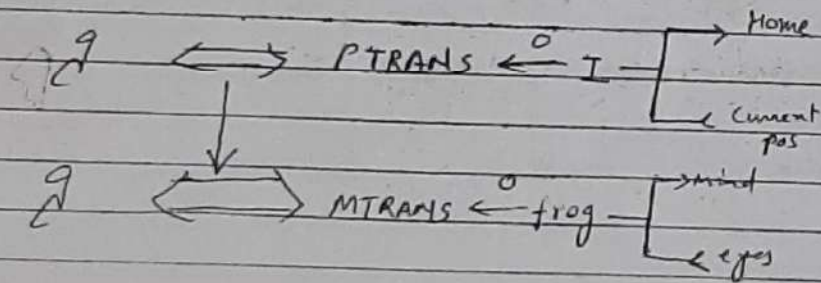
2.

Bill  $\leftrightarrow^P$  INVEST  $\leftarrow^0$  Coke

Q while going home i saw a frog.

As

Q  $\leftrightarrow$  PTRANS  $\leftarrow^0$  frog

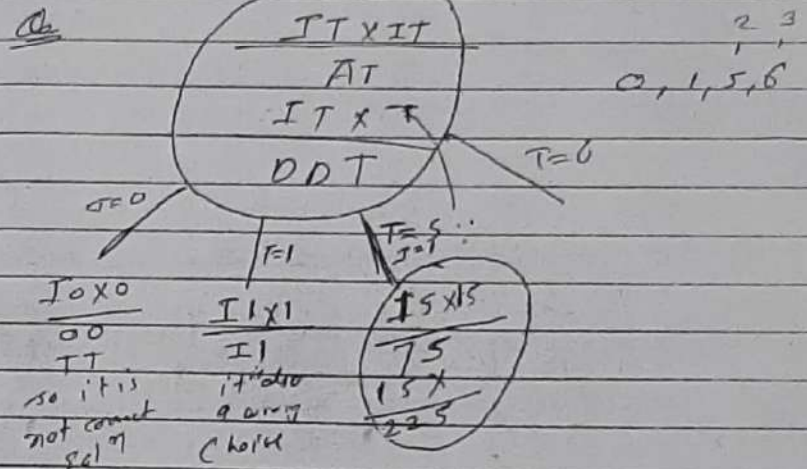


$$\begin{array}{r} ABC \\ ABC \\ \hline CDB \\ BBBX \\ ABCX \\ ABC \end{array}$$

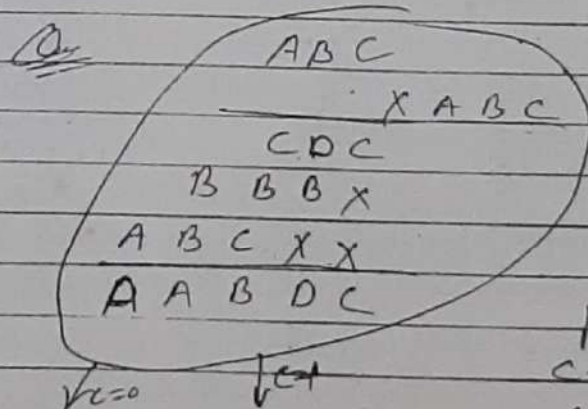
$$\begin{array}{l} A=1 \\ B=0 \end{array} \quad \begin{array}{l} 0,1,6 \\ 2,6 \\ 5,6 \end{array}$$

Constraints Satisfaction Problem  $\rightarrow$   
(Cryptos Arithmetic Problem)

$$\begin{array}{r} SEND \\ + MORE \\ \hline MONEY \end{array} \quad \begin{array}{r} ROADS \\ + CROSS \\ \hline DANNER \end{array}$$



Sol<sup>n</sup> T=5  
 A=7  
 I=1  
 D=2



here we can predict  
 C=5,6  
 B=0

## FUZZY SET THEORY:-

$$A = \{a, b, c, d\}$$

$$A = \{ (a, 0.8), (b, 0.1), (c, 0.0), (d, 1.0) \}$$

(Tall)

↑  
Membership  
char<sup>t</sup> func<sup>n</sup>

$$\text{Short} = \{ (a, 0.1), (b, 0.7), (c, 1.0) \}$$

∴ It may be  
other than 0 & 1  
also

Here probability (Tall) + prob(Short) may not be 1.

Any element of fuzzy set is represent

## FUZZY SET THEORY:-

In conventional set theory only 2-value logic is used i.e. either any element is the member of given set (i.e. probability 1) or it is not member of the set (i.e. prob = 0).

In real life there are large no of app's / situations where partial membership of elements for a particular set is to be represented. The conventional set theory & probability theory does not support such partial membership. Fuzzy theory is the

Sol<sup>n</sup> to handle such type of situation where partial membership can also be considered ex.

ex. Set of tall people & set of short people.

In conventional set theory

$$\text{Tall} = \{a, b, c, d\}$$

$$\text{Short} = \{c, d, g\}$$

i.e. here there two sets are disjoint sets, because it is assumed that either the person is tall or short.

In real life tall is a relative word (not absolute) & the def<sup>n</sup> of the tall varies person-to-person. Hence there may be no of people there are considered both tall short & may be something else in public opinion.

To handle such type of situ<sup>n</sup> the fuzzy repr<sup>n</sup> of such type of situ<sup>n</sup> will be -

$$\text{Tall} = \{ (a, 0.8), (b, 0.3), (c, 0.6), (d, 0.4), (e, 0.1), (f, 0.7) \}$$

$$\text{Short} = \{ (b, 0.3), (c, 0.4), (d, 0.5), (e, 0.2), (f, 0.1) \}$$

As shown in the above sets each element of set is represented by a



pair (

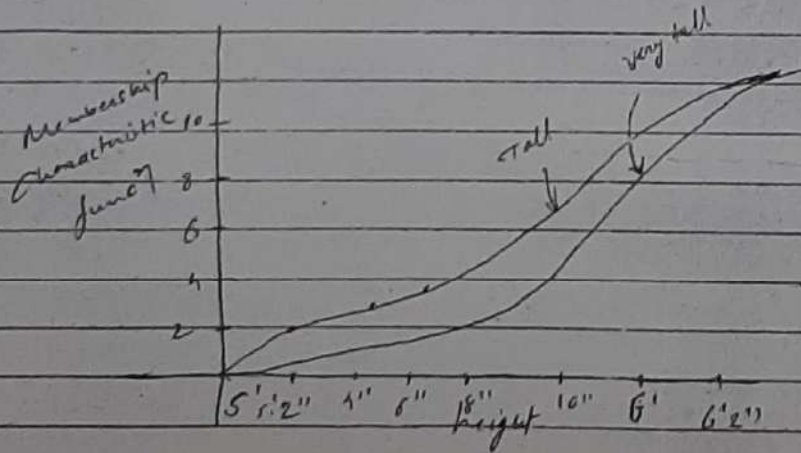
In set tall (B, 0.5) means 50% opinion is that B is tall in set short (B, 0.3) means 30% of people are saying that B is tall remaining 20 people may have not given any opinion or they may consider B as a member of any other set (middle set).

Fuzzy Curve :-

for short, tall & very tall

$$\text{Tall} = \{ (5', 0.0), (5'2'', 0.2), (5'4''), \dots, (6'2'', 1.0) \}$$

Curve



square & for N.V. it will be ...

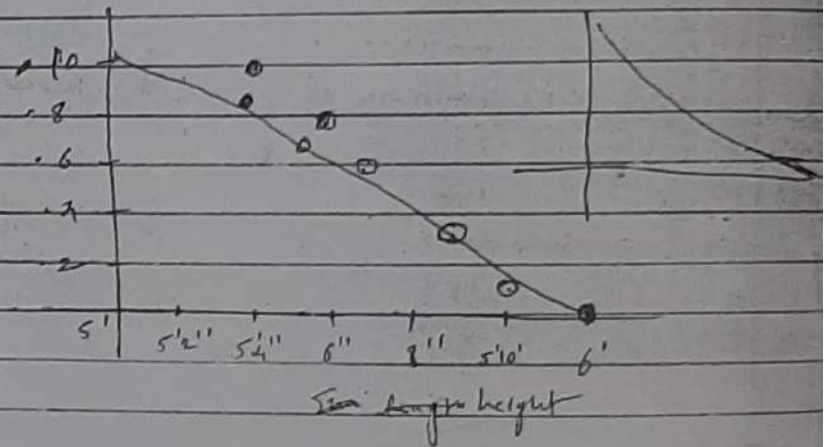
$$\text{Very tall} = \{ (5', 0), (5'2'', 0.04), (5'4'', 0.1), (5'6'', 0.25), \dots, (6', 0.2) \}$$

Or make curve for short & very short (A<sub>2</sub>)

Curve

$$\text{Short} = \{ (3'2'', 1.0), (3'4'', 0.8), (3'6'', 0.7), \dots \}$$

$$\text{Short} = \{ (5', 1.0), (5'2'', 0.8), (5'4'', 0.6), (5'6'', 0.4), (5'8'', 0.2), (6', 0.0) \}$$



## Natural Language Processing →

Converting any natural lang. statement into M/C language instr's (high level)

Phases or For Translation: →

(Steps to translate, N.L. para into M/C lang)

1. Morphological Analysis → It is similar to lexical analysis of compiler. Its basic duty is (a) Consider paragraph as a coll<sup>n</sup> of word  
(b) Verify the validity of each & every word  
(c) For each & every word pass taken & attributes to next phase.

2. Syntactic Analysis: → Here paragraph is considered as coll<sup>n</sup> of sentences and each sentence is verified as per the grammar rules of the nat. lang.

3. Semantic Analysis: - The purpose of this is to check the meaningfulness of each and every sentence.

ex: 1. This is a dog.

2. This dog is

4. Discourse integration: - In natural lang the meaning of a given sentence may be based on previous sentence phrasing  
for ex. (i) Vijay is SATI student  
(ii) He is intelligent.  
(iii) His college is situated at Vidisha.

The purpose of this discourse integration is to get correct meaning of the terms used in the statements.

5. Pragmatic Analysis: → In this it is possible to write phrases and this phrases and sentences don't have any direct meaning but they refer to any other sentence. The purpose of pragmatic analysis is to get actual meaning of the sentence.

CFN For Natural Lang. Process:  $\rightarrow$

$S \rightarrow N_p V_p$  (Noun phrase & Verb phrase)

$N_p \rightarrow \overset{A}{\text{the}} N_p / P_r / N_p / P_n$

$N_p \rightarrow \text{Adj} N$

$\text{Adj} \rightarrow \text{e} / \text{Ad} \text{Adj}$

$V_p \rightarrow V / V N_p$

$N \rightarrow \text{file} / \text{printer} / \text{dog}$

(Proper Noun)  $P_n \rightarrow \text{John} / \text{Bill} / \text{Amit}$

$P_r \rightarrow \text{I}$

$\text{Adj} \rightarrow \text{short} / \text{tall} / \overset{\text{extra}}{\text{good}} / \text{old}$

$V \rightarrow \text{Created} / \text{printed} / \text{bite}$

$A \rightarrow \text{the} / \text{a} / \text{an}$

Q1 Parse the following sentences:

1. Bill printed the file.
2. Sumit is a good boy.

Ans

Q2 The dog <sup>bites</sup> bites the oldman

Ans 1 1. Bill printed the file

Generate by LMD:  $\rightarrow$

$S \rightarrow N_p V_p$

$S \rightarrow P_n V_p$

$S \rightarrow \text{Bill} V_p$

$S \rightarrow \text{Bill} V N_p$

$S \rightarrow \text{Bill printed } N_p$

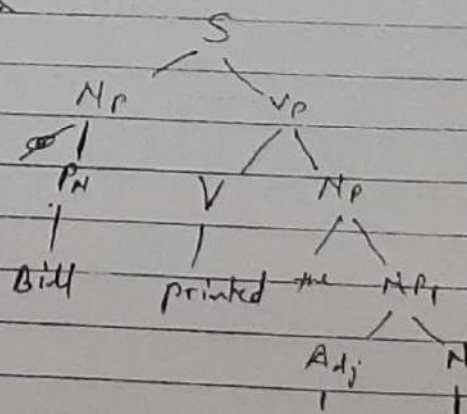
$S \rightarrow \text{Bill printed the } N_p$

$S \rightarrow \text{Bill printed the } \text{Adj} N$

$S \rightarrow \text{---} \text{---} \text{---} \text{the } N$

$S \rightarrow \text{---} \text{---} \text{---} \text{the file}$

parse tree is



S → The dog barks

S → The dog barks NP

S → The dog barks ANP

S → The dog barks AdjN

S → The dog barks AdAdjN

S → The dog barks old N

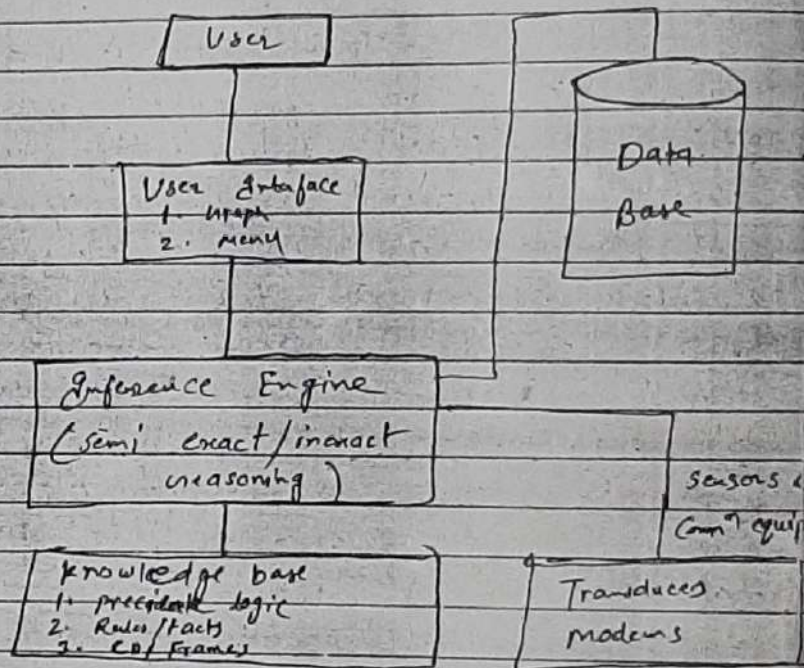
S → The dog barks old man, →

Parse tree →

(Thurs)

## Expert Systems

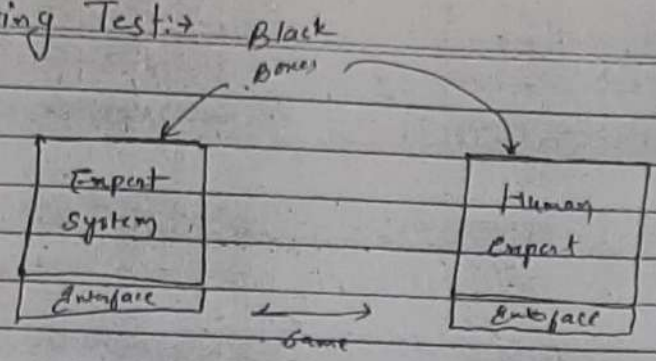
It is generally used by experts. Data is fed to the expert system, which processes the data & do reasoning and then give the by and the results it will give the advice that what to do.



In exact reasoning → means the decision reasoning has not been proved completely but still the decision is taken. Ex. If doctor gets few symptoms of MP parasite & the weather than in spite the proof he takes says it to have the disease.

Cell → The ability to

### Turing Test:



### Expert System Cell:-

A program used to build expert system is usually called expert system cell. It works very similar to word processor which is used to create text file or letters, Expert S.C. is used to create expert system.

### Functions performed by expert system cell:-

- (1) Assists in building the knowledge base by allowing a developer to insert knowledge into knowledge representation structure.
- (2) Provides methods of inference or reduction that can do reasoning based on the info. available in knowledge base.
- (3) Provide an interface that allow user to setup reasoning tasks and query the system about its reasoning strategy.

### What is imp. about E.S. Cell:-

The guidelines for evaluation of expert system cell are as follows:-

- (1) Expressing knowledge:-
  - (a) Descriptive power
  - (b) Ease of translation
  - (c) Readability

### Turing Test:-

The Turing Test is used to verify the goodness of expert systems. It is a black box testing mechanism that is internal components of the expert systems are not tested. Here 2-black boxes are taken in one box human expert is present and in another expert system software. The way of interaction that is interface for both the boxes is same. Now any no. of queries can be asked to the box where we don't know the position of human expert & to expert system soft. and after asking many queries it is not possible to diff. where expert system & where human expert is present then expert system is considered to be perfect and this type of testing mechanism is known as Turing test.

## (2) Organizing & Displaying Knowledge

(a) Reasoning with & validating k

## (3) Reasoning

## (4) Operating a knowledge

(5) Ability to integrate conventional software with expert system.

## Mycin: →

Mycin was one first successful expert system to demonstrate the feasibility of developing intelligible programs that can do activities similar to human experts. Mycin was designed to interact with a physician, collect all the relevant info about the patient and examine this info for the evidence on which a base of diagnosis and recommend a treatment.

Mycin is a rule based system of about 500 rules. In analysing each case mycin answers four basic questions

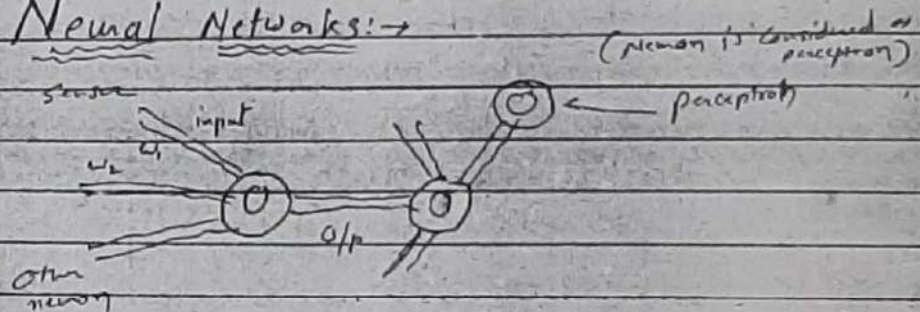
- (1) Is the infection significant.
- (2) What is the identity of the organism causing the infection.
- (3) What are the potentially useful drugs.
- (4) Which drug is best suited for patient

The basic op<sup>n</sup> of Mycin is controlled by backward chaining Rule Interpreter.

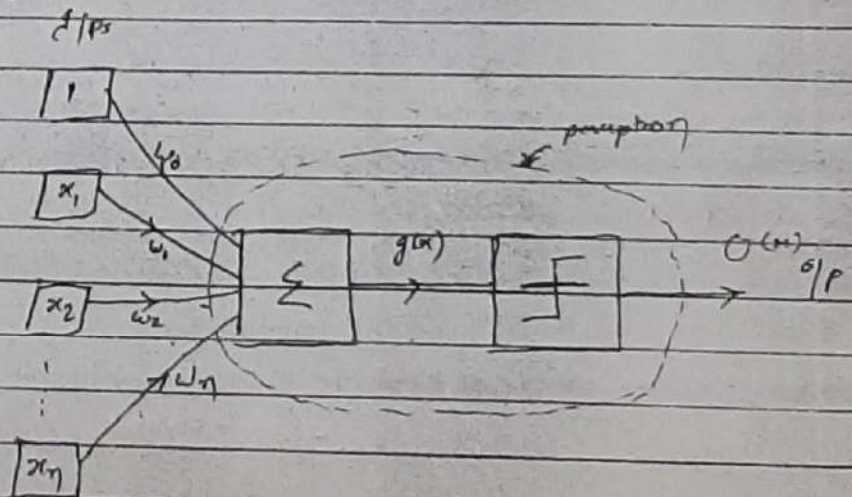
The rules are written with the use of certainty factors. This will help mycin to do - inexact reasoning. The certainty value may range from -1 to +1 here '-1' indicates that complete confidence that the part<sup>l</sup> rule/prop. is false. '+1' indicates the rule - " - true.

(In book)

## Neural Networks: →



## Perceptron: -

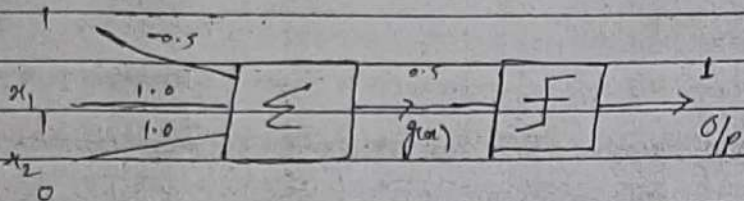


Neural N/w is a complex n/w of perceptrons

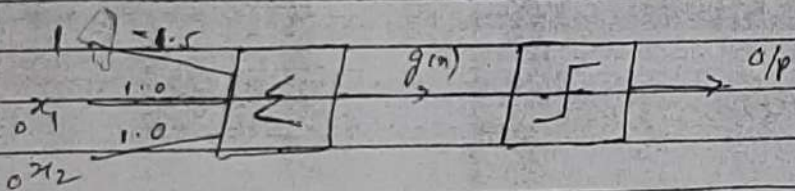
$$g(x) = \sum_{i=0}^n W_i X_i$$

$$O(x) = \begin{cases} 1 & \text{if } g(x) > 0 \\ 0 & \text{if } g(x) < 0 \end{cases}$$

Ex: OR-gate for 2 inp.

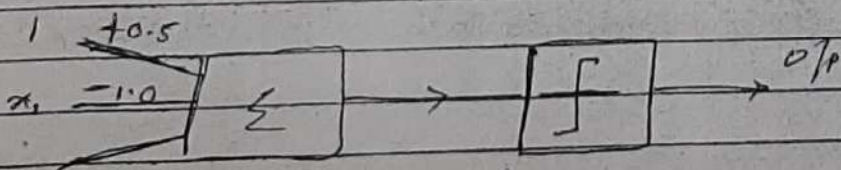


Q AND-gate



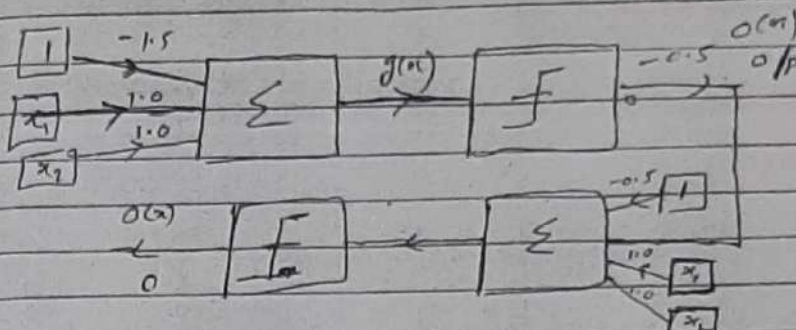
Here

Q EXOR NOT-gate



Q EXOR gate:→

Here we need two pe



LEARNING:-

As human can grow his knowledge base by watching or doing new thing called learning. In AI system learning is provided.

AI system should remain intelligent for long period of time is an essential requirement. To make it possible there must be good knowledge acquisition techniques (Learning tech.) to grow its knowledge base. There are five strategies or learning methods.

- ① Rule learning.
- ② Direct instr<sup>n</sup>
- ③ Analogy.
- ④ Induction
- ⑤ Deduction